

ENDS AND VERTICES OF SMALL DEGREE IN INFINITE MINIMALLY k -(EDGE)-CONNECTED GRAPHS*

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Abstract. Bounds on the minimum degree and on the number of vertices attaining it have been much studied for finite edge-/vertex-minimally k -connected/ k -edge-connected graphs. We give an overview of the results known for finite graphs and show that most of these carry over to infinite graphs if we consider ends of small degree as well as vertices.

Key words. infinite graph, degree, minimum degree, k -connectivity, vertex-minimal, edge-minimal

AMS subject classifications. 05C35, 05C63, 05C40

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1. Introduction.

1.1. The situation in finite graphs. Four notions of minimality will be of interest in this paper. For $k \in \mathbb{N}$, call a graph G *edge-minimally k -connected*, resp., *edge-minimally k -edge-connected*, if G is k -connected, resp., k -edge-connected, but $G - e$ is not for every edge $e \in E(G)$. Analogously call G *vertex-minimally k -connected*, resp., *vertex-minimally k -edge-connected*, if G is k -connected, resp., k -edge-connected, but $G - v$ is not for every vertex $v \in V(G)$. These four classes of graphs often appear in the literature under the names of k -minimal/ k -edge-minimal/ k -critical/ k -edge-critical graphs.

It is known that finite graphs which belong to one of the classes defined above have vertices of small degree. In fact, in three of the four cases, the trivial lower bound of k on the minimum degree is attained. We summarize the known results in the following theorem (some of these results, and similar ones for digraphs, also appear in [1, 9]).

THEOREM 1. *Let G be a finite graph, and let $k \in \mathbb{N}$.*

- (a) (Halin [13]) *If G is edge-minimally k -connected, then G has a vertex of degree k .*
- (b) (Chartrand, Kaugars, and Lick [6], Mader [18]) *If G is vertex-minimally k -connected, then G has a vertex of degree at most $\frac{3}{2}k - 1$.*
- (c) (Lick [17]) *If G is edge-minimally k -edge-connected, then G has a vertex of degree k .*
- (d) (Mader [23]) *If G is vertex-minimally k -edge-connected, then G has a vertex of degree k .*

Note that in Theorem 1(b), the bound of $3k/2 - 1$ on the degree is best possible. For even k , this can be seen by replacing each vertex of C_ℓ , a circle of some length $\ell \geq 4$, with a copy of $K^{k/2}$, the complete graph on $k/2$ vertices, and adding all edges between two copies of $K^{k/2}$ when the corresponding vertices of C_ℓ are adjacent. This procedure is sometimes called the strong product¹ of C_ℓ and $K_{k/2}$. For odd values of

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¹The *strong product* of two graphs H_1 and H_2 is defined in [16] as the graph on $V(H_1) \times V(H_2)$ which has an edge $(u_1, u_2)(v_1, v_2)$ whenever $u_i v_i \in E(H_i)$ for $i = 1$ or $i = 2$ and at the same time either $u_{3-i} = v_{3-i}$ or $u_{3-i} v_{3-i} \in E(H_{3-i})$.

k , similar examples can be constructed, using $K^{(k+1)/2}$'s instead of $K^{k/2}$'s, and in the end deleting two vertices which belong to two adjacent copies of $K^{(k+1)/2}$.

In all four cases of Theorem 1, the minimum degree is attained by more than one vertex.² For convenience let $V_n = V_n(G)$ denote the set of all vertices of a graph G that have degree at most n .

THEOREM 2. *Let G be a finite graph, and let $k \in \mathbb{N}$.*

- (a) (Mader [20]) *In case (a) of Theorem 1, $|V_k| \geq c_k|G|$, where $c_k > 0$ is a constant depending only on k , unless $k = 1$, in which case $|V_k| \geq 2$.*
- (b) (Hamidoune [15]) *In case (b) of Theorem 1, $|V_{3k/2-1}| \geq 2$.*
- (c) (Mader [19, 22]) *In case (c) of Theorem 1, $|V_k| \geq c'_k|G|$, where $c'_k > 0$ is a constant depending only on k , unless $k = 1$ or $k = 3$, in which case $|V_k|$ is at least 2, resp., 4.*
- (d) (Mader [23]) *In case (d) of Theorem 1, $|V_k| \geq 2$.*

In case (a), actually more than the number of vertices of small degree is known: If we delete all the vertices of degree k , we are left with a forest. This was shown in [20]; see also [1]. For extensions of this fact to infinite graphs, see [26].

The difference in the case $k = 1$ in (a) and (c) is due to the paths. For $k = 3$, there is no constant c'_3 in (c): to see this, take the square³ of any long enough path $v_1v_2v_3 \dots v_{\ell-2}v_{\ell-1}v_\ell$, and add the edge v_1v_4 and the edge $v_{\ell-3}v_\ell$. Deleting v_3v_4 and $v_{\ell-3}v_{\ell-2}$, we obtain an edge-minimally 3-edge-connected graph with only six vertices of degree 3.

The constant c_k from (a) can be chosen as $c_k = \frac{k-1}{2k-1}$, and this is best possible [20]. Actually one can ensure [20] that $|V_k| \geq \max\{c_k|G|, k+1, \Delta(G)\}$, where $\Delta(G)$ denotes the maximum degree of G . In (c), the constant c'_k may be chosen as about $1/2$ as well (for estimates, see [2, 5, 24]).

The bounds on the number of vertices of small degree are best possible in (b) and (d) for⁴ $k > 2$. Indeed, for $k \geq 3$, consider the following example. Take any finite number $\ell \geq 2$ of copies H_i of the complete graph $K^{2(k-1)}$, and join every two consecutive H_i with a matching of size $k - 1$ in a way that all these matchings are disjoint. Join a new vertex a to all vertices of H_1 that still have degree $2(k - 1) - 1$, and analogously join a new vertex b to half of the vertices of H_ℓ . Finally join a and b with an edge. See Figure 1.

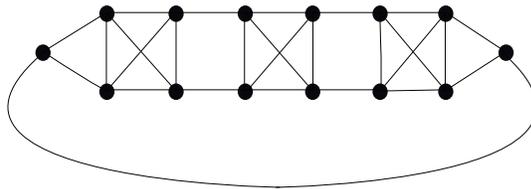


FIG. 1. A finite vertex-minimally k -(edge)-connected graph with only two vertices of degree $< 2(k - 1)$ for $k = 3$.

²We remark that, for uniformity of the results to follow, we do not consider the trivial graph K^1 to be 1-edge-connected/1-connected.

³The square of a graph is obtained by adding an edge between any two vertices of distance 2.

⁴And for $k = 2$, we have $|V_2| \geq 4$ (see [23] for a reference), and this is best possible, as the so-called ladder graphs show. As for $k = 1$, it is easy to see that there are no vertex-minimally 1-(edge)-connected graphs (since we excluded K^1).

The obtained graph is vertex-minimally k -connected as well as vertex-minimally k -edge-connected. However, all vertices but a and b have degree $2(k - 1)$, which, as $k \geq 3$, is greater than $\max\{k, \frac{3}{2}k - 1\}$.

1.2. What happens in infinite graphs? For infinite graphs, a positive result for case (a) of Theorem 1 has been obtained by Halin [14], who showed that every infinite locally finite edge-minimally k -connected graph has infinitely many vertices of degree k , provided that $k \geq 2$. Mader [21] extended the result, showing that, for $k \geq 2$, every infinite edge-minimally k -connected graph G has, in fact, $|G|$ vertices of degree k (see Theorem 3(a) below). It is clear that, for $k = 1$, we are dealing with trees, which, if infinite, need not have any vertices of degree 1.

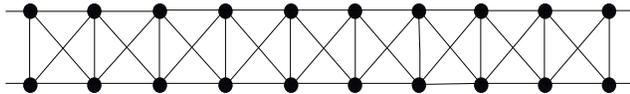


FIG. 2. An infinite vertex-minimally k -connected graph without vertices of degree $3k/2 - 1$ for $k = 2$.

For the other three cases of Theorem 1, the infinite version fails. In fact, for case (b), this can be seen by considering the strong product of the double-ray (i.e., the two-way infinite path) with the complete graph K^k (cf. Figure 2). The obtained graph is $(3k - 1)$ regular and vertex-minimally k -connected. If, instead of the double-ray, we take the r -regular infinite tree T_r for any $r \in \mathbb{N}$, the degrees of the vertices become unbounded in k . For case (d) of Theorem 1, consider the Cartesian product⁵ of K^k with T_r (see Figure 3).

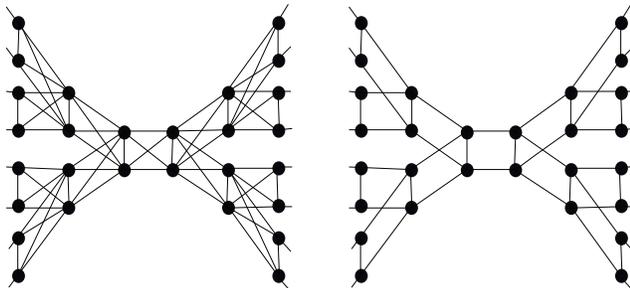


FIG. 3. The strong product (left) and the Cartesian product (right) of T_3 with K^2 .

Counterexamples for an infinite version of (c) will be given now. For the values 1 and 3, this is particularly easy, as for $k = 1$ we may consider the double-ray D and for $k = 3$ its square D^2 . All the vertices of these graphs have degree 2, resp., 4, but D and D^2 are edge-minimally 1-, resp., 3-edge-connected.

For arbitrary values $k \in \mathbb{N}$, we construct a counterexample as follows. Choose $r \in \mathbb{N}$, and take the rk -regular tree T_{rk} . For each vertex v in T_{rk} , insert edges between the neighborhood N_v of v in the next level so that N_v spans r disjoint copies of K^k . (Figure 4 illustrates the case $k = 4, r = 2$.) This procedure gives an edge-minimally k -edge-connected graph, as one easily verifies. However, the vertices of this graph all have degree at least rk .

⁵The Cartesian product of two graphs H_1 and H_2 is defined [7, 16] as the graph on $V(H_1) \times V(H_2)$ which has an edge $(u_1, u_2)(v_1, v_2)$ if, for $i = 1$ or $i = 2$, we have that $u_i v_i \in E(H_i)$ and $u_{3-i} = v_{3-i}$.

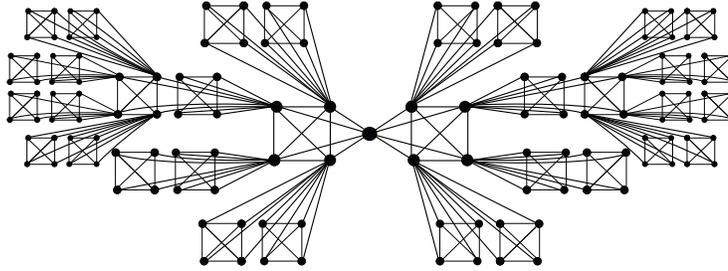


FIG. 4. An edge-minimally 4-edge-connected graph without vertices of degree 4.

Hence a literal extension of Theorems 1 and 2 to infinite graphs is not true except for part (a). The reason can be seen most clearly by comparing Figures 1 and 2: Where in a finite graph we may force vertices of small degree just because the graph has to end somewhere, in an infinite graph we can “escape to infinity.” So an adequate extension of Theorem 1 should also measure something like “the degree at infinity.”

This rather vague-sounding statement can be made precise. In fact, the points “at infinity” are nothing but the ends of graphs, a concept which has been introduced by Freudenthal [10] and later independently by Halin [11] and which is a mainstay of contemporary infinite graph theory. Ends are defined as equivalence classes of rays (one-way infinite paths), where two rays are equivalent if no finite set of vertices separates them. That this is, in fact, an equivalence relation is easy to check. The set of all ends of a graph G is denoted by $\Omega(G)$. For more on ends, consult the infinite chapter of [7].

The concept of the end degree has been introduced in [3] and [25]; see also [7]. In fact, we distinguish between two types of end degrees: the *vertex-degree* and the *edge-degree*. The vertex-degree $d_v(\omega)$ of an end ω is defined as the supremum of the cardinalities of the sets of (vertex)-disjoint rays in ω , and the edge-degree $d_e(\omega)$ of an end ω is defined as the supremum of the cardinalities of the sets of edge-disjoint rays in ω . These suprema are, indeed, maxima [3, 12]. Note that $d_e(\omega) \geq d_v(\omega)$.

In light of this definition, we observe at once what happens in the case $k = 1$ of the infinite version of Theorem 1(a) above. Edge-minimally 1-connected graphs, otherwise known as infinite trees, need not have vertices that are leaves, but if not, then they must have “leaf-like” ends, that is, ends of vertex-degree 1. In fact, it is easy to see that in a tree T with root r , say, every ray starting at r corresponds to an end of T and that all ends of T have vertex- and edge-degree 1. On the other hand, rayless trees have leaves.

This observation gives case (a') in the following generalization of Theorem 1 to infinite graphs. Cases (b)–(d) of Theorem 3, resp., their quantitative versions in Theorem 4, are the main result of this paper.

THEOREM 3. *Let G be a graph, and let $k \in \mathbb{N}$.*

- (a) (Mader [21]) *If G is edge-minimally k -connected and $k \geq 2$, then G has a vertex of degree k .*
- (a') *If G is edge-minimally 1-connected, then G has a vertex of degree 1 or an end of edge-degree 1.*
- (b) *If G is vertex-minimally k -connected, then G has a vertex of degree $\leq \frac{3}{2}k - 1$ or an end of vertex-degree $\leq k$.*

- (c) If G is edge-minimally k -edge-connected, then G has a vertex of degree k or an end of edge-degree $\leq k$.
- (d) If G is vertex-minimally k -edge-connected, then G has a vertex of degree k or an end of vertex-degree $\leq k$.

As in the finite case, one can give bounds on the number of vertices/ends of small degree. Recall that $V_n = V_n(G)$ denotes the set of all vertices of degree at most n , and let $\Omega_n^v = \Omega_n^v(G)$, resp., $\Omega_n^e = \Omega_n^e(G)$, denote the set of ends of vertex-, resp., edge-degree at most n .

THEOREM 4. Let G be a graph, and let $k \in \mathbb{N}$.

- (a) (Mader [21]) In case (a) of Theorem 3, $|V_k| = |G|$,
- (a') In case (a') of Theorem 3, $|V_1 \cup \Omega_1^e| = |G|$ unless $|G| \leq \aleph_0$, in which case $|V_1 \cup \Omega_1^e| \geq 2$.
- (b) In case (b) of Theorem 3, $|V_{3k/2-1} \cup \Omega_k^v| \geq 2$.
- (c) In case (c) of Theorem 3, $|V_k \cup \Omega_k^e| \geq 2$.
- (d) In case (d) of Theorem 3, $|V_k \cup \Omega_k^v| \geq 2$.

Concerning part (c), we remark that we may replace graphs with multigraphs (see Corollary 12). Also, in (a') and (c), one may replace the edge-degree with the vertex-degree as this yields a weaker statement.

We shall prove Theorem 4(b), (c), and (d) in sections 2, 3, and 4, resp. Statement (a') is fairly simple; in fact, it follows from our remark above that every tree has at least two leaves/ends of vertex-degree 1. In general, this is already the best bound because of the (finite or infinite) paths. For trees T of uncountable order, we get more as these have to contain a vertex of degree $|G|$, and it is then easy to find $|G|$ vertices/ends of (edge)-degree 1.

In analogy to the finite case, the bounds on the degrees of the vertices in (b) cannot be lowered, even if we allow the ends to have larger vertex-degree. An example for this is given at the end of section 2. There we also state a lemma that says that the vertex-/edge-degree of the ends in Theorem 4 will, in general, not be less than k .

Also the bound on the number of vertices/ends of small degree in Theorem 4(b) and (d) is best possible. For (d), this can be seen by considering the Cartesian product of the double ray with the complete graph K^k (for $k = 2$, that is, the double-ladder). For (b), we may again consider the strong product of the double-ray with the complete graph K^k (see Figure 2 for $k = 2$). The latter example also shows that in (b), we cannot replace the vertex-degree with the edge-degree.

As for Theorem 4(c), it might be possible that the bound of Theorem 2(c) extends. For infinite graphs G , the positive proportion of the vertices there should translate to an infinite set S of vertices and ends of small degree/edge-degree. More precisely one would wish for a set S of cardinality $|V(G)|$ or, even stronger, $|S| = |V(G) \cup \Omega(G)|$. Observe that it is necessary to exclude also in the infinite case the two exceptional values $k = 1$ and $k = 3$ as there are graphs (e.g., D and D^2) with only two vertices/ends of (edge)-degree 1, resp., 3.

QUESTION 5. For $k \neq 1, 3$, does every infinite edge-minimally k -edge-connected graph G contain infinitely many vertices or ends of (edge)-degree k ? Does G have $|V(G)|$, or even $|V(G) \cup \Omega(G)|$, such vertices or ends?

Another interesting question is which k -(edge)-connected graphs have vertex- or edge-minimally k -(edge)-connected subgraphs. Especially interesting in the case of edge-connectivity would be an edge-minimally k -edge-connected subgraph on the same vertex set as the original graph. Finite graphs trivially do have such subgraphs, but for infinite graphs this is not always true. One example in this respect is the double-ladder,

which is 2-connected but has no edge-minimally 2-connected subgraphs on the same vertex set. This observation leads to the study of vertex-/edge-minimally k -(edge)-connected (*standard*) *subspaces* rather than graphs. For more on this, see [8, 26], the latter of which contains a version of Theorem 3(a) for standard subspaces.

We finish the introduction with a few necessary definitions. The *vertex-boundary* $\partial_v H$ of a subgraph H of a graph G is the set of all vertices of H that have neighbors in $G - H$. The *edge-boundary* of H is the set $\partial_e H = E(H, G - H)$. A *region* of a graph is a connected induced subgraph H with finite vertex-boundary $\partial_v H$. If $\partial_v H = k$, then we call H a k -*region* of G . A *profound region* is a region H with $H - \partial_v H \neq \emptyset$.

2. Vertex-minimally k -connected graphs. In this section we shall show part (b) of Theorem 4. For the proof, we need two lemmas. The first of these lemmas may be extracted⁶ from [6] or from [21] and at once implies Theorem 1(b). For completeness, we shall give a proof.

LEMMA 6. *Let $k \in \mathbb{N}$, let G be a vertex-minimally k -connected graph, and let H be a profound finite k -region of G . Then G has a vertex v of degree at most $\frac{3}{2}k - 1$. Moreover, if $|G - H| > |H - \partial_v H|$, then we may choose $v \in V(H)$.*

Proof. Note that we may assume H is inclusion-minimal with the above properties. Set $T := \partial_v H$, set $C_1 := H - T$, and set $C_2 := G - H$. Let $x \in V(C_1)$, and observe that since G is vertex-minimally k -connected, there is a k -separator T' of G with $x \in T'$. Let D_1 be a component of $G - T'$, set $D_2 := G - T' - D_1$, and set $T^* := T \cap T'$. Furthermore, for $i, j = 1, 2$, set $A_j^i := C_i \cap D_j$ and set $T_j^i := (T' \cap C_i) \cup (T \cap D_j) \cup T^*$. Observe that $N(A_j^i) \subseteq T_j^i$.

We claim that there are i_1, i_2, j_1, j_2 with either $(i_1, j_1) = (i_2, 3 - j_2)$ or $(i_1, j_1) = (3 - i_2, j_2)$ such that, for $(i, j) \in \{(i_1, j_1), (i_2, j_2)\}$,

$$(2.1) \quad |T_j^i| \leq k \text{ and } A_j^i = \emptyset.$$

In fact, observe that, for $j = 1, 2$, we have that $|T_j^1| + |T_{3-j}^2| = |T| + |T'| = 2k$. Thus either $|T_j^1| \leq k$, which by the minimality of H implies that A_j^1 is empty, or $|T_{3-j}^2| < k$, which by the k -connectivity of G implies that A_{3-j}^2 is empty. This proves (2.1).

We hence know that there is an $X \in \{C_1, C_2, D_1, D_2\}$ such that $V(X) \subseteq T \cup T'$. Now, as $|T| = |T'| = k$, (2.1) implies that

$$2|X| + k + |T^*| \leq |T_{j_1}^{i_1}| + |T_{j_2}^{i_2}| \leq 2k$$

and hence that

$$(2.2) \quad |X| + \frac{|T^*|}{2} \leq \frac{k}{2}.$$

Thus there is a vertex $v \in X$ of degree at most

$$\max\{|T| + |X| - 1, |T'| + |X| - 1\} \leq k + k/2 - 1.$$

Clearly we may assume $v \in V(H)$ unless both $|T_1^1|$ and $|T_2^1|$ are strictly greater than k . But then by (2.1), $V(C_2) \subseteq T'$, and thus by (2.2), $|C_2| \leq k/2 - |T^*|/2$. So

$$|C_2| \leq k - |T^*| - |C_2| = |T'| - |T^*| - |C_2| \leq |T' \cap C_1| \leq |C_1|$$

as desired. \square

⁶Although the graphs there are all finite, the procedure is the same.

We also need a lemma from [25].

LEMMA 7 (see [25, Lemma 5.2]). *Let G be a graph such that all its ends have vertex-degree at least $m \in \mathbb{N}$. Let C be an infinite region of G . Then there exists a profound region $C' \subseteq C$ for which one of the following holds:*

- (a) C' is finite and $|\partial_v C'| < m$, or
- (b) C' is infinite and $|\partial_v C''| \geq m$ for every profound region $C'' \subsetneq C'$.

Observe that the outcome of Lemma 7 is invariant under modifications of the structure of $G - C$. Hence, in all applications, we may assume that $d_v(\omega) \geq m$ only for ends ω of G that have rays in C .

We are now ready to prove Theorem 4(b).

Proof of Theorem 4(b). First of all, we claim that, for every infinite region H of G , it holds that

$$(2.3) \quad \text{there is a vertex } v \in V(H) \text{ of degree } \leq \frac{3}{2}k - 1 \text{ or an end of vertex-degree } \leq k \text{ with rays in } H.$$

In order to see (2.3), we assume that there is no end as desired and apply Lemma 7 to H with $m := k + 1$. This yields a profound region $H' \subseteq H$. We claim that (a) of Lemma 7 holds; then we may use Lemma 6 to find a vertex $w \in V(H')$ with $d(w) \leq 3k/2 - 1$.

So assume for contradiction that (b) of Lemma 7 holds. Since G is k -connected, there exists a finite family \mathcal{P} of finite paths in G such that each pair of vertices from $\partial_v H'$ is connected by k pairwise internally disjoint paths from \mathcal{P} . Set

$$S := \partial_v H' \cup V\left(\bigcup \mathcal{P}\right),$$

and observe that $H' - S$ is still infinite. In particular, $H' - S$ contains a vertex v .

Since G is vertex-minimally k -connected, v lies in a k -separator T' of G . By the choice of $v \notin S$, no two vertices of $\partial_v H'$ are separated by T' . Thus all of $\partial_v H' - T'$ is contained in one component of $G - T'$.

Let C'' be a component of $G - T'$ that does not contain any vertices from $\partial_v H'$. Note that as G is k -connected, v has a neighbor in C'' . Hence $C'' \subseteq H' - \partial_v H'$, and $H'' := G[C'' \cup T']$ is a profound region with $H'' \subseteq H'$.

In fact, $H'' \subsetneq H'$, which is clear if $H' = G$ and otherwise follows from the fact that $v \in T' - \partial_v H'$, and thus, because $|\partial_v H'| \geq k = |T'|$, we know that $\partial_v H' - T' \neq \emptyset$. So (b) implies that $k + 1 \leq |T'| = k$, a contradiction as desired. This proves (2.3).⁷

Now let $T \subseteq V(G)$ be any separator of G of size k (such a separator exists by the vertex-minimality of G). First suppose that $G - T$ has at least one infinite component C . Then we apply Lemma 6 or (2.3) to any component of $G - C$ and find an end ω of vertex-degree k with no rays in C or a vertex $v \in V(G - C)$ of degree at most $3k/2 - 1$. Let x denote the point found, that is, $x = \omega$ or $x = v$.

Let C' be the subgraph of G induced by C and all vertices of T that have infinite degree into C . Then C' is a region, and we may thus apply (2.3) to C' in order to find the second end/vertex of small (vertex)-degree. This second point is different from x by the choice of C' .

It remains to treat the case when all components of $G - T$ are finite. As we otherwise apply Theorem 2(b), we may assume that $G - T$ has infinitely many components. Hence, as G has no $(k - 1)$ -separators, each $x \in T$ has infinite degree. This means that we may apply Lemma 6 to any two k -regions H_1, H_2 with $\partial_v H_1 = T = \partial_v H_2$

⁷Observe that taking $H = G$, we have thus proved Theorem 3(b).

in order to find two vertices $v_1 \in V(H_1) - \partial_v H_1$, $v_2 \in V(H_2) - \partial_v H_2$, each of degree $\leq 3k/2 - 1$. \square

We remark that the bound on the vertex-degree given by Theorems 3(b) and 4(b) is best possible. Indeed, by the following lemma, which follows from Lemma 7.1 from [4], the vertex-degree of the ends of a k -connected locally finite graph has to be at least k .

LEMMA 8. *Let $k \in \mathbb{N}$, let G be a locally finite graph, and let $\omega \in \Omega(G)$. Then $d_v(\omega) = k$ if and only if k is the smallest integer such that every finite set $S \subseteq V(G)$ can be separated⁸ from ω with a set of k vertices.*

For nonlocally finite graphs, the minimum size of an S - ω separator corresponds to the vertex-/edge-degree of ω plus the number of dominating vertices of ω . See [4].

One might now ask whether it is possible to achieve a better upper bound on the degree of the “small degree vertices” than the one given by Theorems 3 and 4(b) if one accepts a worse bound on the vertex-degree of the “small degree ends.” The answer is no. This is illustrated by the following example (cf. Figure 5) for even k . For odd k , there are similar examples.

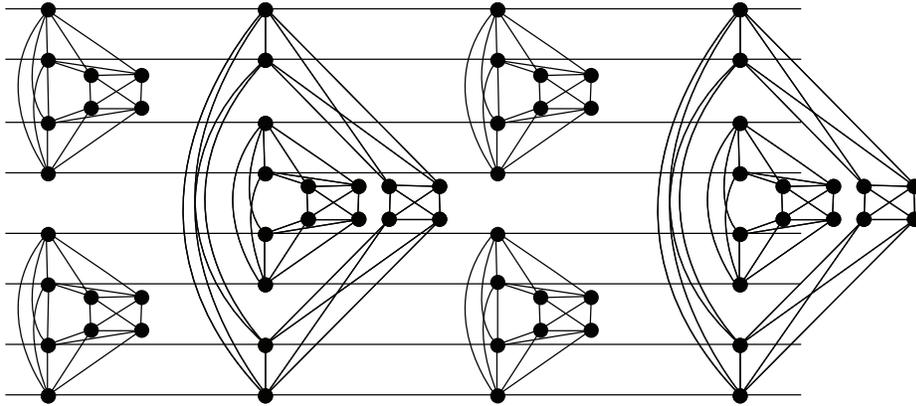


FIG. 5. A vertex-minimally k -connected graph G with $d(v) \geq \frac{3}{2}k - 1$ and $d_v(\omega) \gg k$ for all $v \in V(G)$ and $\omega \in \Omega(G)$ for $k = 4$.

Let $\ell \in \mathbb{N} \cup \{\aleph_0\}$, and take the disjoint union of ℓ double-rays R_1, \dots, R_ℓ . For simplicity, assume that k divides ℓ . For each $i \in \mathbb{Z}$, take ℓ/k copies of the strong product of C_4 with $K^{k/2}$, and identify the vertices that belong to the first or the last copy of $K^{k/2}$ with the i th vertices the R_j . This can be done in a way so that the obtained graph, which is easily seen to be vertex-minimally k -connected, has two ends of vertex-degree ℓ , while the vertices have degree either $3k/2 - 1$ or $3k/2 + 1$.

3. Edge-minimally k -edge-connected graphs. We now prove part (c) of Theorem 4. We start by proving a lemma that will be useful also in section 4.

LEMMA 9. *Let G be a graph, and let $(D_i)_{i \in \mathbb{N}}$ be a sequence of regions such that $D_{i+1} \subseteq D_i - \partial_v D_i$ for all $i \in \mathbb{N}$. Then there is an end $\omega \in \Omega(G)$ that has a ray in each of the D_i so that*

- (i) if $|\partial_v D_i| \leq k$ for all i , then $d_v(\omega) \leq k$, and

⁸We say a set $T \subseteq V(G)$ separates a set $S \subseteq V(G)$ from an end $\omega \in \Omega(G)$ if the unique component of $G - T$ that contains rays of ω does not contain vertices from S .

(ii) if $|\partial_e D_i| \leq k$ for all i , then $d_e(\omega) \leq k$.

Proof. As all the D_i are connected, it is easy to construct a ray R which has a subray in each of the D_i . Say R belongs to the end $\omega \in \Omega(G)$. We show only (i) as (ii) can be proved analogously.

Suppose, for contradiction, that $|\partial_v D_i| \leq k$ for all i , but $d_v(\omega) > k$. Then ω contains a set \mathcal{R} of $k + 1$ disjoint rays. Let S be the set of all starting vertices of these rays. Since $D_i \subseteq D_{i-1} - \partial_v D_{i-1}$ for all i , there is an $n \in \mathbb{N}$ such that $S \cap V(D_n) = \emptyset$. (To be precise, one may take $n := \max_{s \in S, v \in \partial_v D_0} \text{dist}(s, v) + 1$.) But then it is impossible that all rays of \mathcal{R} have subrays in D_n as only k of them can pass disjointly through $\partial_v D_n$. \square

We also need the following lemma from [25].

LEMMA 10 (see [25, Lemma 3.2]). *Let $m \in \mathbb{N}$, and let $D \neq \emptyset$ be a region of a graph G so that $|\partial_e D| < m$ and so that $|\partial_e D'| \geq m$ for every nonempty region $D' \subseteq D - \partial_v D$ of G . Then there is an inclusion-minimal nonempty region $H \subseteq D$ with $|\partial_e H| < m$.*

Combined, the two lemmas yield a lemma similar to Lemma 7 from the previous section.

LEMMA 11. *Let $D \neq \emptyset$ be a region of a graph G so that $|\partial_e D| < m$ and so that $d_e(\omega) \geq m$ for every end $\omega \in \Omega(G)$ with rays in D . Then there is an inclusion-minimal nonempty region $H \subseteq D$ with $|\partial_e H| < m$.*

Proof. Set $D_0 := D$, and, inductively for $i \geq 1$, choose a nonempty region $D_i \subseteq D_{i-1} - \partial_v D_{i-1}$ such that $|\partial_e D_i| < m$ (if such a region D_i exists). If at some step i we are unable to find a region D_i as above, then we apply Lemma 10 to D_{i-1} to find the desired region H . On the other hand, if we end up defining an infinite sequence of regions, then Lemma 9(ii) tells us that there is an end ω with rays in D and that $d_e(\omega) < m$, a contradiction. \square

We are now ready to prove part (c) of our main theorem.

Proof of Theorem 4(c). Since G is edge-minimally k -edge-connected, G has a nonempty region D such that $|\partial_e D| = k$ and such that $G - D \neq \emptyset$. We shall find a vertex or end of small (edge)-degree in D ; then one may repeat the procedure for $G - D$ in order to find the second point.

First we apply Lemma 11 with $m := k + 1$ to obtain an end as desired or an inclusion-minimal nonempty region $H \subseteq D$ with $|\partial_e H| \leq k$. If $V(H)$ should consist of only one vertex, then this vertex has degree k as desired. So suppose that $V(H)$ has more than one vertex, that is, $E(H)$ is not empty.

Let $e \in E(H)$. By the edge-minimal k -edge-connectivity of G , we know that e belongs to some cut F of G with $|F| = k$. Say, $F = E(A, B)$, where $A, B \neq \emptyset$ partition $V(G)$. Since $e \in F$, neither $A_H := A \cap V(H)$ nor $B_H := B \cap V(H)$ is empty.

So $|\partial_e A_H| > k$ and $|\partial_e B_H| > k$ by the minimality of H . But then, since $|\partial_e H| \leq k$ and $|F| \leq k$, we obtain that

$$\begin{aligned} |\partial_e(A \setminus A_H)| + |\partial_e(B \setminus B_H)| &\leq 2|\partial_e H| + 2|F| - |\partial_e A_H| - |\partial_e B_H| \\ &< 4k - 2k \\ &= 2k. \end{aligned}$$

Hence at least one of $|\partial_e(A \setminus A_H)|$, $|\partial_e(B \setminus B_H)|$, say, the former, is strictly smaller than k . Since G is k -edge-connected, this implies that $A \setminus A_H$ is empty. But then $A \subsetneq V(H)$, a contradiction to the minimality of H . \square

We dedicate the rest of this section to multigraphs, that is, graphs with parallel edges, which often appear to be the more appropriate objects when studying

edge-connectivity. Defining the edge-degree of an end ω of a multigraph in the same way as for graphs, that is, as the supremum of the cardinalities of the sets of edge-disjoint rays from ω , and defining V_k and Ω_k^e as earlier for graphs, we may apply the proof of Theorem 4(c) with only small modifications⁹ to multigraphs. We thus get the following corollary.

COROLLARY 12. *Let G be an edge-minimally k -edge-connected multigraph. Then $|V_k(G) \cup \Omega_k(G)^e| \geq 2$.*

In particular, every finite edge-minimally k -edge-connected multigraph has at least two vertices of degree k .

However, a statement in the spirit of Theorem 2(c) does not hold for multigraphs, no matter whether they are finite or not. For this, it suffices to consider the graph we obtain by multiplying the edges of a finite or infinite path by k . This operation results in an edge-minimally k -edge-connected multigraph which has no more than the two vertices/ends of (edge)-degree k which were promised by Corollary 12.

4. Vertex-minimally k -edge-connected graphs. In this section we prove Theorem 4(d). The proof is based on Lemma 14, which at once yields Theorem 2(d), the finite version of Theorem 4(d). The idea of the proof of this lemma (in particular, Lemma 13) is similar to Mader’s original proof of Theorem 2(d) in [23].¹⁰

We need one auxiliary lemma before we get to Lemma 14. For a set $X \subseteq V(G) \cup E(G)$ in a graph G , write $X_V := X \cap V(G)$ and $X_E := X \cap E(G)$.

LEMMA 13. *Let $k \in \mathbb{N}$. Let G be a graph, let $S \subseteq V(G) \cup E(G)$ with $|S| \leq k$, and let C be a component of $G - S$ so that $|C| \leq |S_E|$. Then C contains a vertex of degree at most k .*

Proof. Suppose that the vertices of C all have degree at least $k + 1$. Then each sends at least $k + 1 - |S_V| - (|C| - 1)$ edges to $G - S - C$. This means that

$$|C|(k + 1 - |S_V| - (|C| - 1)) \leq |S_E| \leq k - |S_V|.$$

So $|C|(k - |S_V| - |C| + 1) \leq k - |S_V| - |C|$, which, as $|C| \geq 1$, is possible only if both sides of the inequality are negative, that is, if $|C| > k - |S_V|$. But this is impossible as $|C| \leq |S_E| \leq k - |S_V|$. \square

As usual, the edge-connectivity of a graph G is denoted by $\lambda(G)$. Also, in order to make clear which underlying graph we are referring to, it will be useful to write $\partial_e^G H = \partial_e H$ for a region H of a graph G .

LEMMA 14. *Let $k \in \mathbb{N}$, let G be a k -edge-connected graph, and let C be an inclusion-minimal region of G with the property that C has a vertex x so that $|\partial_e^{G-x}(C - x)| = \lambda(G - x) < k$ and $C - x \neq \emptyset$. Suppose, for each $y \in V(C)$, the graph $G - y$ has a cut of size less than k . Then $C - x$ contains a vertex of degree k (in G).*

Proof. If every vertex of $C - x$ has a neighbor in $D := G - x - C$, then we may apply Lemma 13 with $S := \{x\} \cup \partial_e^{G-x}(C - x)$ and are done. So let us assume that there is a vertex $y \in V(C - x)$, all of whose neighbors lie in C . By assumption, $G - y$ has a cut F of size $\lambda(G - y) < k$, which splits $G - y$ into A and B , with $x \in V(A)$, say. See Figure 6.

⁹We will then have to use a version of Lemma 10 for multigraphs. Observe that such a version holds as we may apply Lemma 10 to the (simple) graph obtained by subdividing all edges of the multigraph. This procedure will not affect the degrees of the ends. The rest of the proof will then go through replacing everywhere “graph” with “multigraph.”

¹⁰But as [23] does not contain the statement we need for finding the second vertex/end of small degree, we cannot make use of it here.

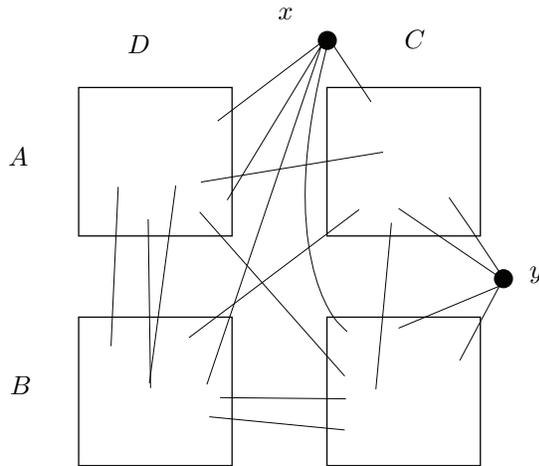


FIG. 6. The graph G with $A, B, C, D, x,$ and y .

Since G is k -edge-connected, F is not a cut of G . Hence y has neighbors in both A and B . Thus, as $N(y) \subseteq V(C)$ and $x \in V(A)$, it follows that $B \cap C \neq \emptyset$. Consider the region C' induced by $B \cap C$ and y . Because $x \in V(A)$, we know that $C' \subsetneq C$.

So, by the choice of C and x , we may assume that $|\partial_e^{G-y}(B \cap C)| > \lambda(G-y) = |F|$. Thus

$$\begin{aligned} |\partial_e^{G-x}(A \cap D)| &\leq |\partial_e^{G-x}(C-x)| + |F| - |\partial_e^{G-y}(B \cap C)| \\ &< |\partial_e^{G-x}(C-x)| \\ &= \lambda(G-x), \end{aligned}$$

implying that $A \cap D = \emptyset$. That is, $A \cup y \subsetneq C$. (Here we use again that $B \cap C \neq \emptyset$.) As $|F| = \lambda(G-y) < k$, this contradicts the minimality of C . \square

Any finite vertex-minimally k -edge-connected graph G clearly has an inclusion-minimal region C as in Lemma 14. Thus Theorem 1(d) follows at once from Lemma 14. Applying Lemma 14 to any inclusion-minimal region with the desired properties that is contained in $G - (C - x)$ in order to find a second vertex of small degree in G , we get the following corollary.

COROLLARY 15 (see Theorem 2(d)). *Let G be a finite vertex-minimally k -edge-connected graph. Then G has at least two vertices of degree k .*

This means that, for a proof of Theorem 4(d), we need to worry about only the infinite regions, which is accomplished in the next lemma.

LEMMA 16. *Let $k \in \mathbb{N}$, let G be a vertex-minimally k -edge-connected graph, and let D be a region of G . Let $x \in V(D)$ such that $|\partial_e^{G-x}(D-x)| = \lambda(G-x) < k$. Suppose G has no inclusion-minimal region $C \subseteq D$ with the property that C contains a vertex y so that $|\partial_e^{G-y}(C-y)| = \lambda(G-y) < k$. Then G has an end of vertex-degree $\leq k$ with rays in D .*

Proof. We construct a sequence of infinite regions D_i of G , starting with $D_0 := D$ which clearly is infinite. Our regions will have the property that $D_i \subseteq D_{i-1} - \partial_v D_{i-1}$, which means that we may apply Lemma 9(i) in order to find an end as desired.

In step $i \geq 1$, for each pair of vertices in $\partial_v^G D_{i-1}$, take a set of k edge-disjoint paths joining them: The union of all these paths gives a finite subgraph H of G . Since D_{i-1} was infinite, $D_{i-1} - H$ still is and thus contains a vertex y .

Since G is vertex-minimally k -edge-connected, $G - y$ has a cut of size less than k , which splits $G - y$ into A and B , say, which we may assume to be connected. Say, A contains a vertex of $\partial_v^G D_{i-1}$. Then $\partial_v^G D_{i-1} \subseteq V(A)$ (since $y \notin V(H)$). Thus, as y has neighbors in both A and B (because G is k -connected), we obtain that $B \subseteq D_{i-1}$. Observe that $D_i := B \cup y$ is infinite as otherwise it would contain an inclusion-minimal region C as prohibited in the statement of the lemma. \square

We finally prove Theorem 4(d).

Proof of Theorem 4(d). Let $x \in V(G)$, and let F be a cut of $G - x$ with $|F| = \lambda(G - x) < k$. Say, F splits $G - x$ into A_1 and A_2 . For $i = 1, 2$, if A_i contains an inclusion-minimal region C such that C has a vertex y with the property that $|\partial_e^{G-y}(C - y)| = \lambda(G - y) < k$, we use Lemma 14 to find a vertex of degree at most k in $C - y \subseteq A_i$. On the other hand, if A_i does not contain such a region, we use Lemma 16 to find an end of the desired vertex-degree. \square

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