

# On the Erdős-Menger conjecture with ends

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## Abstract

A well-known conjecture of Erdős states that, given an infinite graph  $G$  and sets  $A, B \subseteq V(G)$ , there exists a family of disjoint  $A$ - $B$  paths  $\mathcal{P}$  together with an  $A$ - $B$  separator  $X$  consisting of a choice of one vertex from each path in  $\mathcal{P}$ . There is a natural extension of this conjecture in which  $A$ ,  $B$  and  $X$  may contain ends as well as vertices. We prove this extension for sets  $A$  and  $B$  that can be separated by countably many vertices or ends, and for sets  $A$  and  $B$  which have disjoint closures in the end topology of  $G$ .

## 1 Introduction

Erdős conjectured (see [11]) that Menger's theorem should extend to infinite graphs as follows:

**Erdős-Menger Conjecture.** *For every graph  $G = (V, E)$  and any two sets  $A, B \subseteq V$  there is a set  $\mathcal{P}$  of disjoint  $A$ - $B$  paths in  $G$  and an  $A$ - $B$  separator  $X$  consisting of a choice of one vertex from each of the paths in  $\mathcal{P}$ .*

There are several partial results [3]. In particular, Aharoni [2] proved the conjecture for countable graphs. In [4], this was extended as follows:

**Theorem 1.1.** [4] *The Erdős-Menger conjecture holds for all graphs  $G = (V, E)$  and sets  $A, B \subseteq V$  that are separated in  $G$  by a countable set of vertices.*

In particular, the conjecture holds whenever  $A$  is countable, regardless of the cardinality of  $G$ .

Another line of attack was taken in [7]. Again, an assumption is made that  $A$  and  $B$  are easy to separate. But this time, the notion of separation used is topological:

**Theorem 1.2.** [7] *The Erdős-Menger conjecture holds for all graphs  $G = (V, E)$  and sets  $A, B \subseteq V$  whose closures in the topological space  $|G|$  consisting of  $G$  together with its ends are disjoint.*

(See Section 2 for a formal definition of  $|G|$ , and some discussion of what the disjoint closures condition means for the relative position of  $A$  and  $B$  in  $G$ .)

Although Theorem 1.2 refers implicitly to the ends of  $G$  by its closure condition, the conclusion is the original one from Erdős's conjecture, which makes no reference to ends. However, there is also a natural extension of the conjecture that does refer to ends. Here, the sets  $A$  and  $B$  may contain ends as well as vertices. The  $A$ - $B$  paths in  $\mathcal{P}$  can be either finite paths linking two vertices,

or rays linking a vertex to an end, or double rays linking two ends. Similarly, the separator  $X$  may contain ends (that lie in  $A$  or  $B$ ), thus blocking any ray belonging (= converging) to that end. This extension was proposed in [6], and found to be true for countable graphs under certain necessary restrictions for  $A$  and  $B$ .

In this paper, we extend both Theorems 1.1 and 1.2 to ends, in the spirit of [6]. Our extension of Theorem 1.2 will build on our extension of Theorem 1.1.

## 2 Definitions and statement of results

The basic terminology we use can be found in [5]. All the graphs in this paper are simple and undirected. As most of the graphs we deal with are infinite, we recall a few (standard) concepts for infinite graphs. Let  $G = (V, E)$  be a fixed infinite graph.

A 1-way infinite path is called a *ray*, a 2-way infinite path a *double ray*. The subrays of rays or double rays are their *tails*. The *ends* of  $G$  are the equivalence classes of rays under the following equivalence relation: two rays  $R_1, R_2$  in  $G$  are *equivalent* if no finite set of vertices separates them. As one easily observes, this condition holds if and only if there are infinitely many disjoint (finite)  $R_1$ - $R_2$  paths. This in turn is equivalent to the existence of a ray that meets both  $R_1$  and  $R_2$  infinitely often. The set of ends of  $G$  is denoted by  $\Omega = \Omega(G)$ , and we write  $G = (V, E, \Omega)$  to refer to  $G$  together with its set of ends.

*Paths* in  $G$  can be finite paths (which contain at least one vertex), rays, double rays, or singleton sets  $\{\omega\}$ , where  $\omega$  is an end of  $G$ .

We will now define the standard topology on  $G$  together with its ends. We start by viewing  $G$  itself as a 1-complex. (Thus, the basic open neighbourhoods of an inner point of an edge are the open intervals on the edge containing that point, while the basic open neighbourhoods of a vertex  $v$  are the unions of half-open intervals containing  $v$ , one from every edge at  $v$ .<sup>1</sup> The point set of this 1-complex will again be denoted by  $G$ .) To extend this topology to the set  $G \cup \Omega$ , we have to define a neighbourhood basis for every end  $\omega \in \Omega$ . To do so, consider any finite set  $S \subseteq V$ . Then  $G - S$  has exactly one component  $C = C(S, \omega)$  that contains a tail of every ray in  $\omega$ . We say that  $\omega$  *belongs* to  $C$ , and write  $\overline{C}(S, \omega)$  for the component  $C$  together with all the ends of  $G$  belonging to  $C$ . As the basic open neighbourhoods of  $\omega$  we now take all sets of the form

$$\hat{C}(S, \omega) := \overline{C}(S, \omega) \cup E'(S, \omega),$$

where  $S$  is any finite subset of  $V$  and  $E'(S, \omega)$  is any union of half-edges  $(z, y] \subseteq e$ , one for every edge  $e = xy$  with  $x \in S$  and  $y \in C$  (with  $z \in e$ ).

Let  $|G|$  denote the topological space on  $G \cup \Omega$  thus defined. (When  $G$  is locally finite,  $|G|$  is a compact space known as the *Freudenthal compactification* of  $G$ ; see [8, 9] for more.)

We write  $\overline{X}$  for the closure of a set  $X \subseteq |G|$  in  $|G|$ . For example, the set  $\overline{C}(S, \omega)$  defined above is the closure in  $|G|$  of the set  $C(S, \omega)$ . Generally, the

<sup>1</sup>Alternatively, we might fix for every edge  $[u, v]$  a homeomorphism with the real interval  $[0, 1]$  and take as basic open neighbourhoods for a vertex  $v$  only those unions of half-open edges  $[v, z)$  whose images in  $[0, 1]$  have the same length. This gives a different topology when vertices have infinite degrees, but since all relevant sequences of points will be sequences of vertices or of ends, the difference does not matter.

difference between a subgraph  $H$  and its closure  $\overline{H}$  is always a set of ends of  $G$  (possibly empty). These need not correspond to ends of  $H$  and should not be confused with them. For example, if  $G$  is the 1-way infinite ladder and  $H$  consists of all the rungs, then  $\overline{H} \setminus H$  consists of one point, the unique end  $\omega$  of  $G$ . But  $H$  itself has no ends. Similarly, the subgraph  $H' = G - E(H)$  of  $G$  is a double ray and thus has two ends, but  $\overline{H'} \setminus H' = \{\omega\}$  as before.

With precise definitions now available, let us take another look at what the assumption of  $\overline{A} \cap \overline{B} = \emptyset$  in Theorem 1.2 means for the relative position of the sets  $A$  and  $B$  (which we assume to be disjoint). Formally, the condition means that every ray in  $G$  can be separated by finitely many vertices from at least one of the sets  $A$  and  $B$ . An obvious way to ensure this, of course, is to assume that some finite set of vertices separates  $A$  from  $B$ . But this is much stronger (except when  $G$  is locally finite), and the conjecture has long been known for this case. A more typical example of  $\overline{A} \cap \overline{B} = \emptyset$  is to take as  $A$  and  $B$  two distinct levels of vertices in the  $\kappa$ -regular tree (for  $\kappa$  any infinite cardinal), in which case  $G$  contains  $\kappa$  disjoint  $A$ - $B$  paths.

The closure of an infinite path  $P$  contains one or two ends of  $G$ . (Even if  $P$  is a double ray, its closure may contain only one end, as in the ladder example above.) We will often consider such an end as the first or last point of  $P$ , and when we say that two paths are *disjoint* then these points too shall be distinct. (The first and last point of a path  $P = \{\omega\}$ , of course, is  $\omega$ .) For  $A, B \subseteq V \cup \Omega$ , a path is an  $A$ - $B$  *path* if its first but no other point lies in  $A$  and its last but no other point lies in  $B$ .

A set  $X \subseteq V \cup \Omega$  is an  $A$ - $B$  *separator* in a subspace  $T \subseteq |G|$  if every path  $P$  in  $T$  with its first point in  $A$  and its last point in  $B$  satisfies  $P \cap X \neq \emptyset$ . (We express this informally by saying that “ $P$  meets  $X$ ”, though strictly speaking we shall mean  $\overline{P}$  rather than just  $P$ .) We say that a set  $Y \subseteq V \cup \Omega$  lies *on* a set  $\mathcal{P}$  of disjoint  $A$ - $B$  paths if  $Y$  consists of a choice of exactly one vertex or end from every path in  $\mathcal{P}$ . We say that  $G$  *satisfies the Erdős-Menger conjecture for  $A$  and  $B$* , or that *the Erdős-Menger conjecture holds for  $G, A, B$* , if  $|G|$  contains a set  $\mathcal{P}$  of disjoint  $A$ - $B$  paths and an  $A$ - $B$  separator on  $\mathcal{P}$ . (Thus, officially, we always refer to the ends version of the conjecture. But this is compatible with the traditional terminology: if neither  $A$  nor  $B$  contains an end then neither can any  $A$ - $B$  path, so the conjecture with ends automatically defaults to the original conjecture in this case.)

The union of a ray  $R$  and infinitely many disjoint paths starting on  $R$  but otherwise disjoint from  $R$  is a *comb* with *spine*  $R$ . The last points (vertices or ends) of those paths are the *teeth* of the comb. We will frequently use the following simple lemma:

**Lemma 2.1.** *In the graph  $G = (V, E, \Omega)$  let  $R$  be a ray of an end  $\omega$ , and let  $X \subseteq V \cup \Omega$  such that  $\omega \notin X$ . Then  $\omega \in \overline{X}$  if and only if  $G$  contains a comb with spine  $R$  and teeth in  $X$ .  $\square$*

We can now state our two main results. First, our extension of Theorem 1.1:

**Theorem 2.2.** *Let  $G = (V, E, \Omega)$  be a graph, let  $A, B \subseteq V \cup \Omega$  satisfy  $A \cap \overline{B} = \emptyset = \overline{A} \cap B$ , and suppose that there exists a countable  $A$ - $B$  separator  $X \subseteq V \cup \Omega$  in  $G$ . Then  $G$  satisfies the Erdős-Menger conjecture for  $A$  and  $B$ .*

In particular, the ends version of the Erdős-Menger conjecture is true whenever  $A$  is countable. We remark that the condition  $A \cap \overline{B} = \emptyset = \overline{A} \cap B$  is really

necessary [6]; so Theorem 2.2 is best possible in this respect. If neither  $A$  nor  $B$  contains an end, the condition reduces to  $A \cap B = \emptyset$ , which we may always assume without loss of generality; thus, Theorem 2.2 implies Theorem 1.1.

Our second main result, whose proof builds on Theorem 2.2, extends Theorem 1.2 fully to its natural topological setting:

**Theorem 2.3.** *Every graph  $G = (V, E, \Omega)$  satisfies the Erdős-Menger conjecture for all sets  $A, B \subseteq V \cup \Omega$  that have disjoint closures in  $|G|$ .*

Let us complete this section with an outline of the proofs to come, and of how the paper is organized.

The proof of Theorem 2.2 will occupy us for the next two sections. It runs roughly as follows. Most of the proof – all of Section 3 – will be spent on transferring our problem to an equivalent problem in which  $G$  is replaced with a suitable minor and  $A$  and  $B$  consist of vertices only. In this new situation, the countable  $A$ – $B$  separator  $X$  – which likewise may be assumed to consist of vertices only – divides  $G$  into two parts: one between  $X$  and  $A$  (including both) and the rest (which includes  $B$ ). We now apply Theorem 1.1 to obtain an  $A$ – $X$  separator  $Y$  on a system of disjoint  $A$ – $X$  paths in the first part. Note that  $Y$  is again countable, and it separates  $A$  from  $B$  in  $G$ . Repeating the same procedure for the part of  $G$  between  $Y$  and  $B$  yields a system of  $Y$ – $B$  paths with a separator  $Z$  on it, which again separates  $A$  from  $B$ . These paths can be concatenated with the  $A$ – $Y$  segments of the first, to give a system of disjoint  $A$ – $B$  paths (with the  $A$ – $B$  separator  $Z$  on it). It remains to transfer this solution back to the original sets  $A, B$  containing ends, in the original graph  $G$ .

In Section 5 we prove Theorem 2.3. Employing techniques developed in Section 3 and in [7], we will eliminate all ends in the closures of  $A$  and  $B$ . Then the remaining ends can be discarded as well. In this way, the problem is reduced to a rayless graph, for which the Erdős-Menger conjecture is known to hold:

**Theorem 2.4** (Aharoni [1], Polat [12]). *The Erdős-Menger conjecture holds for rayless graphs.*

### 3 The reduction lemma

In this section we develop further some techniques from [6] designed to reduce the ends versions of the Erdős-Menger conjecture to the related vertex versions. Observe that in the finite Menger theorem we can ignore all the vertices in  $A \cap B$  and work with the graph  $G - (A \cap B)$  instead. In an infinite graph, however, we have to take care that no end in  $A \cup B$  is destroyed or split when the vertices of  $A \cap B$  are deleted from  $G$ .

**Lemma 3.1.** *Let  $G = (V, E, \Omega)$  be a graph, and let  $A, B \subseteq V \cup \Omega$  satisfy*

$$A \cap (\overline{B} \setminus B) = \emptyset = (\overline{A} \setminus A) \cap B.$$

*Then for the graph  $G' := G - (A \cap B \cap V)$  there are sets  $A', B' \subseteq V(G') \cup \Omega(G')$  satisfying*

- (i)  $|A'| \leq |A|$ ;
- (ii) if  $A \subseteq V$  then  $A' \subseteq A$ , and if  $B \subseteq V$  then  $B' \subseteq B$ ;

(iii)  $A' \cap \overline{B'} = \emptyset = \overline{A'} \cap B'$ ;

(iv) if  $G'$  satisfies the Erdős-Menger conjecture for  $A'$  and  $B'$ , then  $G$  satisfies it for  $A$  and  $B$ .

*Proof.* Put  $A' := A \setminus B$  and  $B' := B \setminus A$ , both of which are subsets of  $|G|$ . Consider a ray  $R$  of an end  $\alpha$  in  $A'$  or  $B'$ , say in  $A'$ . Then  $R$  has a tail in  $G'$ . Indeed, if not then there are vertices of  $A \cap B \cap V \subseteq B$  in every neighbourhood of  $\alpha \in A \setminus B$ . Consequently,  $\alpha \in A \cap (\overline{B} \setminus B)$ , which is a contradiction. Similarly, two rays  $R_1, R_2$  in  $G'$  of which  $R_1$  is a ray of an end  $\omega \in A' \cup B'$  are equivalent in  $G'$  if and only if they are equivalent in  $G$ . Indeed, if  $R_1$  and  $R_2$  are equivalent in  $G$  then there is a ray  $R_3 \in \omega$  that meets both of  $R_1$  and  $R_2$  infinitely often. Now  $R_3$  has a tail in  $G'$ , showing that  $R_1$  and  $R_2$  are also equivalent in  $G'$ .

Thus, mapping every end of  $G$  in  $A' \cup B'$  to the unique end of  $G'$  that contains tails of its rays defines a bijection between the ends in  $A' \cup B'$  and certain ends in  $G'$ . Using this bijection (and a slight abuse of notation) we may view  $A'$  and  $B'$  also as subsets of  $V(G') \cup \Omega(G')$ . Clearly,  $A' \cap B'$  is empty and hence (iii) is satisfied. Also, (i) and (ii) are trivial.

For (iv), let  $X'$  be an  $A'$ - $B'$  separator on a set of disjoint  $A'$ - $B'$  paths  $\mathcal{P}'$  in  $G'$ . Adding to  $\mathcal{P}'$  the trivial paths  $\{x\}$  for all  $x \in A \cap B$  yields a set  $\mathcal{P}$  of disjoint  $A$ - $B$  paths on the  $A$ - $B$  separator  $X := X' \cup (A \cap B)$ .  $\square$

Later on, in Lemma 3.5, we shall need a family of disjoint subgraphs of  $G$  (with certain properties) such that every end of  $A$  lies in the closure of one of these subgraphs. Such a family cannot always be found. But our next lemma finds instead a family of subgraphs such that the ends of  $A$  not contained in their closures form a set  $I$  that can be ignored: those ends will automatically be separated from  $B$  by any  $(A \setminus I)$ - $B$  separator on a set of disjoint  $A$ - $B$  paths.

**Lemma 3.2.** *Let  $G = (V, E, \Omega)$  be a graph, and let  $A, B \subseteq V \cup \Omega$  be such that  $A \cap \overline{B} = \emptyset = \overline{A} \cap B$ . Then for every set  $A_\Omega \subseteq A \cap \Omega$  there exist a set  $I \subseteq A_\Omega$ , an ordinal  $\mu^*$ , and families  $(G_\mu)_{\mu < \mu^*}$  and  $(S_\mu)_{\mu < \mu^*}$  such that, for every  $\mu < \mu^*$ , the graph  $G_\mu - S_\mu$  is a component of  $G - S_\mu$  with  $S_\mu$  as its finite set of neighbours, and*

(i)  $\overline{G_\mu - S_\mu} \cap B = \emptyset$ ;

(ii) if  $G_\mu \neq \emptyset$  then  $\overline{G_\mu} \cap A_\Omega \neq \emptyset$ ;

(iii)  $V(G_\nu \cap G_\mu) \subseteq S_\nu \cap S_\mu$  for all  $\nu < \mu$ .

Moreover,

(iv) for every end  $\alpha \in A_\Omega \setminus I$  there is a  $\mu < \mu^*$  with  $\alpha \in \overline{G_\mu}$ ;

(v) every  $(A \setminus I)$ - $B$  separator on a set of disjoint  $(A \setminus I)$ - $B$  paths is also an  $A$ - $B$  separator.

*Proof.* We construct the families  $(G_\mu)_{\mu < \mu^*}$  and  $(S_\mu)_{\mu < \mu^*}$  and a transfinite sequence  $I_0 \subseteq I_1 \subseteq \dots \subseteq A_\Omega$  recursively. The sets  $I_\mu$  ( $\mu < \mu^*$ ) will serve as precursors to  $I$ . To simplify notation, we write  $C_\mu := G_\mu - S_\mu$  for every  $\mu$ . For the construction, we will in addition to (i)–(iii) require for every  $\mu$  that

(vi)  $I_\mu \cap \overline{G_\nu} = \emptyset$  for all  $\nu \leq \mu$ .

We start by setting  $I_0, G_0, S_0 := \emptyset$ . Consider the least ordinal  $\mu > 0$  such that the above sets are already defined for all  $\lambda < \mu$ . If  $\mu$  is a limit, we set

$$I_\mu := \bigcup_{\lambda < \mu} I_\lambda$$

and  $G_\mu, S_\mu := \emptyset$ . This choice clearly satisfies (i)–(iii) and (vi).

Suppose now that  $\mu$  is a successor,  $\mu = \lambda + 1$  say. If every end in  $A_\Omega \setminus I_\lambda$  lies in some  $\overline{G_\nu}$  with  $\nu < \mu$ , we set  $\mu^* := \mu$  and terminate the recursion. So suppose there is an end  $\alpha \in A_\Omega \setminus I_\lambda$  that lies in no earlier  $\overline{G_\nu}$ . Then, if possible, choose a finite vertex set  $S$  such that  $C(S, \alpha)$  avoids all  $G_\nu$  with  $\nu < \mu$ .

Such a choice of  $S$  is impossible if and only if

$$\text{for every finite } S \subseteq V \text{ there is a } \nu < \mu \text{ with } C(S, \alpha) \cap G_\nu \neq \emptyset. \quad (1)$$

In this case we choose to ignore  $\alpha$ , i.e. set  $I_\mu := I_\lambda \cup \{\alpha\}$  and  $G_\mu, S_\mu := \emptyset$ . Again the requirements (i)–(iii) are clearly met, while (vi) holds by the choice of  $\alpha$ .

Now suppose we can find  $S$  as desired. As  $A \cap \overline{B} = \emptyset$ , we can also find a basic open neighbourhood  $\hat{C}(S', \alpha)$  of  $\alpha$  in  $|G|$  that is disjoint from  $B$ . We now define  $S_\mu$  as the set of neighbours of  $C(S \cup S', \alpha)$  and  $G_\mu := G[S_\mu \cup C(S_\mu, \alpha)]$ . Then (i) holds since  $S_\mu \supseteq S'$ , while (ii) holds as  $\alpha \in \overline{G_\mu}$ . To see (iii), first note that

$$G_\nu \cap C_\mu = \emptyset \text{ for all } \nu < \mu$$

by the choice of  $S$ . So, all we have to show is that  $G_\nu \cap S_\mu \subseteq S_\nu$ . Consider a vertex  $v \in G_\nu \cap S_\mu$ . Since  $S_\mu$  is the set of neighbours of  $C_\mu$ , there is a vertex  $w \in C_\mu$  adjacent to  $v$ . As noted above,  $w \notin G_\nu$ . So  $v$  is a vertex in  $G_\nu = C_\nu \cup N(C_\nu)$  with a neighbour outside  $G_\nu$ , implying  $v \notin C_\nu$  and hence  $v \in S_\nu$ , as desired.

Let us finally set  $I_\mu := I_\lambda$  and verify (vi). We only need to show that  $I_\mu \cap \overline{G_\mu} = \emptyset$ . Suppose that intersection contains an end  $\alpha'$ . Let  $\mu' < \mu$  be minimal such that  $\alpha' \in I_{\mu'}$ . Then (1) should have been satisfied for  $\mu'$  and  $\alpha'$ , but fails with  $S := S_\mu$  as  $C(S_\mu, \alpha') = C_\mu$ , a contradiction.

Having defined  $I_\mu, G_\mu$  and  $S_\mu$  for all  $\mu < \mu^*$  so that (i)–(iii) and (vi) are satisfied, we put

$$I := \bigcup_{\mu < \mu^*} I_\mu.$$

Together with the definition of  $\mu^*$  this implies (iv). Observe that from (vi) we obtain  $I \cap \overline{G_\mu} = \emptyset$  for all  $\mu < \mu^*$ .

To establish (v) let  $\mathcal{P}$  be a system of disjoint  $(A \setminus I)$ – $B$  paths and  $X$  an  $(A \setminus I)$ – $B$  separator on  $\mathcal{P}$ . Now suppose that  $X$  is not an  $A$ – $B$  separator in  $|G|$ , i.e. there is a path  $Q$  from  $A$  to  $B$  that avoids  $X$ . By turning  $Q$  into a path  $\tilde{Q}$  from  $A \setminus I$  to  $B$  that avoids  $X$ , we will obtain a contradiction.

We may assume that  $\tilde{Q}$  starts at an end  $\alpha \in I$ . Let  $\mu$  be the step at which  $\alpha$  was added to  $I$ , i.e. let  $\mu$  be minimal with  $\alpha \in I_\mu$ . Choose a finite vertex set  $S$  such that  $\overline{C}(S, \alpha)$  is disjoint from  $B$  (this is possible, as  $A \cap \overline{B} = \emptyset$ ). Then any path of  $\mathcal{P}$  that meets  $C(S, \alpha)$  must pass through  $S$ . Hence only finitely many paths of  $\mathcal{P}$  can meet  $C(S, \alpha)$ , and so  $X_\alpha := X \cap \overline{C}(S, \alpha)$  is also finite. Conditions (iii) and (iv) ensure that every end in  $X_\alpha$  lies in exactly one  $\overline{C_\lambda}$ ; let  $\{\lambda_1, \dots, \lambda_m\}$  be the set of these  $\lambda$ . Then for

$$S' := S \cup (X_\alpha \cap V) \cup \bigcup_{i=0}^m S_{\lambda_i}$$

we have

$$\overline{C}(S', \alpha) \cap X = \emptyset.$$

Now, all we need is a point of  $A \setminus I$  that lies in  $\overline{C}(S', \alpha)$  (and thus can be used to change  $Q$  into the desired path). Indeed, if there is an ordinal  $\lambda < \mu$  such that  $G_\lambda \neq \emptyset$  and

$$C_\lambda \subseteq C(S', \alpha), \quad (2)$$

we can complete the proof as follows. By (ii) for  $\lambda$  there will be an end  $\alpha' \in A$  in  $\overline{C}_\lambda \subseteq \overline{C}(S', \alpha)$ . Since  $I \cap \overline{G}_\lambda = \emptyset$ , we have  $\alpha' \in A \setminus I$ . Take an  $\alpha'$ - $Q$  path  $P$  in  $\overline{C}(S', \alpha)$  with last vertex  $x$ , say. Then  $P$  avoids  $X$ , and hence so does the path  $\tilde{Q} := PxQ$ . Thus,  $\tilde{Q}$  is as desired.

So suppose there is no ordinal  $\lambda < \mu$  satisfying (2). Then for all  $\lambda < \mu$  we have either  $C_\lambda \cap C(S', \alpha) = \emptyset$  or  $C_\lambda \cap S' \neq \emptyset$ . As all the  $C_\lambda$  are disjoint by (iii), only finitely many of them meet  $S'$ ; let  $\lambda_{m+1}, \dots, \lambda_n$  be the corresponding ordinals. Then

$$S'' := S' \cup \bigcup_{i=m+1}^n S_{\lambda_i}$$

satisfies  $C(S'', \alpha) \cap C_\lambda = \emptyset$  for all  $\lambda < \mu$ .

However,  $G_\lambda \cap C(S'', \alpha)$  cannot be empty for all  $\lambda < \mu$ , as this would contradict (1) for step  $\mu$  with  $S := S''$ . So there exists an ordinal  $\lambda < \mu$  with  $S_\lambda \cap C(S'', \alpha) \neq \emptyset$ . A vertex  $v$  in this intersection must have a neighbour in  $C_\lambda$ , which then also lies in  $S' \cup C(S', \alpha)$  because  $C(S'', \alpha) \subseteq C(S', \alpha)$ . Thus,

$$(S' \cup C(S', \alpha)) \cap C_\lambda \neq \emptyset.$$

Since  $C_\lambda \not\subseteq C(S', \alpha)$  by assumption, this implies that  $C_\lambda$  meets  $S'$ . But then  $\lambda \in \{\lambda_{m+1}, \dots, \lambda_n\}$  and hence  $S_\lambda \subseteq S''$ , contradicting the fact that  $v$  lies in both  $S_\lambda$  and  $C(S'', \alpha)$ .  $\square$

For our end-to-vertex reduction we need two more lemmas.

**Lemma 3.3.** [6] *Let  $H$  be a subgraph of a graph  $G$ , let  $S \subseteq V(H)$  be finite, and let  $T \subseteq V(H) \cup \Omega(G)$  be such that  $T \subseteq \overline{H}$ . Then  $\overline{H}$  contains a set  $\mathcal{P}$  of disjoint  $S$ - $T$ -paths and an  $S$ - $T$ -separator (in  $\overline{H}$ ) on  $\mathcal{P}$ .*

For a set  $T$  of vertices in a graph  $H$ , a  $T$ -path is a path that meets  $T$  only in its first and last vertex. A set of paths will be called *disjoint outside* a given subgraph  $Q \subseteq H$  if distinct paths meet only in  $Q$ .

**Lemma 3.4.** [13, 6] *Let  $H$  be a graph,  $T \subseteq V(H)$  finite, and  $k \in \mathbb{N}$ . Then  $H$  has a subgraph  $H'$  containing  $T$  such that for every  $T$ -path  $Q = s \dots t$  in  $H$  meeting  $H - H'$  there are  $k$  distinct  $T$ -paths from  $s$  to  $t$  in  $H'$  that are disjoint outside  $Q$ .*

Our next lemma allows us to replace the set  $A \subseteq V \cup \Omega$  in Theorem 2.2 with a set  $A'$  consisting only of vertices.

**Lemma 3.5.** *Let  $G = (V, E, \Omega)$  be a graph, and let  $A, B \subseteq V \cup \Omega$  be such that  $A \cap B = \emptyset = \overline{A} \cap B$ . Then there are a minor  $G' = (V', E', \Omega')$  of  $G$  and sets  $A' \subseteq V', B' \subseteq V' \cup \Omega'$  satisfying*

- (i)  $|A'| \leq |A|$ ;

- (ii) if  $B \subseteq V$  then  $B' \subseteq B$ ;
- (iii)  $A' \cap \overline{B'} = \emptyset = \overline{A'} \cap B'$ ;
- (iv)  $G$  satisfies the Erdős-Menger-conjecture for  $A$  and  $B$  if  $G'$  satisfies it for  $A'$  and  $B'$ .

*Proof.* Applying Lemma 3.2 with  $A_\Omega := A \cap \Omega$  we obtain an ordinal  $\mu^*$ , subgraphs  $G_\mu$ , finite vertex sets  $S_\mu$  and a set of ends  $I \subseteq A$ . Our aim is to change  $G$  into  $G'$  by deleting and contracting certain connected subgraphs of our graphs  $G_\mu - S_\mu$ . By Lemma 3.2 (iii) we shall be able to do this independently for the various  $G_\mu$ : for each  $\mu < \mu^*$  separately, we shall find in  $G_\mu - S_\mu$  a set  $\mathcal{D}_1(\mu)$  of connected subgraphs to be deleted, and another set  $\mathcal{D}_2(\mu)$  of connected subgraphs that will be contracted.

Fix  $\mu < \mu^*$ . If  $G_\mu$  is empty we let  $\mathcal{D}_1(\mu) = \mathcal{D}_2(\mu) = \emptyset$ . Assume now that  $G_\mu \neq \emptyset$ . Put  $A_\mu := A \cap \overline{G_\mu}$ . Applying Lemma 3.3 to  $H = G_\mu$  we find in  $\overline{G_\mu}$  a finite set  $\mathcal{P}$  of disjoint  $S_\mu$ - $A_\mu$  paths and an  $S_\mu$ - $A_\mu$  separator  $X_\mu$  on  $\mathcal{P}$ . We write  $X_\mu = U_\mu \cup O_\mu$ , where  $U_\mu = X_\mu \cap V$  and  $O_\mu = X_\mu \cap \Omega$ , both of which are finite since  $|X_\mu| \leq |\mathcal{P}| \leq |S_\mu|$ . Moreover,

$$U_\mu \text{ separates } S_\mu \text{ from } A_\mu \setminus O_\mu \text{ in } G. \quad (3)$$

Indeed, every  $S_\mu$ - $(A_\mu \setminus O_\mu)$  path in  $G$  lies in  $\overline{G_\mu}$  and hence meets  $X_\mu$ , and since it cannot meet  $O_\mu$  unless it ends there, it meets  $X_\mu$  in  $U_\mu$ .

We define  $\mathcal{D}_1(\mu)$  as the set of all the components  $D$  of  $G - U_\mu$  whose closure  $\overline{D}$  meets  $A_\mu \setminus O_\mu$ . By (3), these components satisfy  $D \subseteq G_\mu - S_\mu$ , and their neighbourhood  $N(D) \subseteq U_\mu$  in  $G$  is finite. In addition,

$$\overline{D} \cap O_\mu = \emptyset \text{ for all } D \in \mathcal{D}_1(\mu). \quad (4)$$

For if  $\alpha \in \overline{D} \cap O_\mu$ , say, and  $P$  is the  $S_\mu$ - $A_\mu$  path in  $\mathcal{P}$  that ends in  $\alpha$ , then  $P$  has a tail in  $D$ . Since  $P$  does not meet  $U_\mu \supseteq N(D)$ , this implies  $P \subseteq \overline{D}$ . Consequently,  $S_\mu \cap D$  is not empty as it contains at least the first vertex of  $P$ . This contradicts  $D \subseteq G_\mu - S_\mu$ .

Put

$$H_\mu := G_\mu - \bigcup \mathcal{D}_1(\mu).$$

Note that, as every  $v \in U_\mu$  lies on a path in  $\mathcal{P}$ ,

$$G_\mu \text{ contains a set of disjoint } H_\mu\text{-}A_\mu \text{ paths whose set of first points is } U_\mu. \quad (5)$$

By (3) and the definition of  $H_\mu$ , we have  $\overline{H_\mu} \cap A \subseteq U_\mu \cup O_\mu = X_\mu$ . Since  $O_\mu$  is finite, we can extend  $U_\mu \cup S_\mu$  to a finite set  $T_\mu \subseteq V(H_\mu)$  that separates the ends in  $O_\mu$  pairwise in  $G$ . Let  $H'_\mu$  be the finite subgraph of  $H_\mu$  containing  $T_\mu$  which Lemma 3.4 provides for  $k := |S_\mu| + 1$ , and for each  $\alpha \in O_\mu$  let  $D_\alpha$  be the component of  $G - H'_\mu$  to which  $\alpha$  belongs. Finally, we conclude our definitions for  $\mu$  by setting  $\mathcal{D}_2(\mu) := \{D_\alpha \mid \alpha \in O_\mu\}$ .

Define for  $i = 1, 2$

$$\mathcal{D}_i := \bigcup_{\mu < \mu^*} \mathcal{D}_i(\mu).$$

Observe that, by Lemma 3.2 (iii) and since their neighbourhoods in  $G$  are finite, the elements of  $\mathcal{D}_1 \cup \mathcal{D}_2$  have pairwise disjoint closures.

Before we can define  $G'$ , we first have to introduce a graph  $\tilde{G} = (\tilde{V}, \tilde{E}, \tilde{\Omega})$  from which we will obtain  $G'$  by deleting certain vertices. Let  $\tilde{G}$  be obtained from  $G - \bigcup \mathcal{D}_1$  by contracting every  $D_\alpha \in \mathcal{D}_2$  to a single vertex  $a_\alpha$ , and put

$$A^* := \{a_\alpha \mid D_\alpha \in \mathcal{D}_2\}.$$

Then for  $Z := \bigcup \mathcal{D}_1 \cup \bigcup \mathcal{D}_2$  we have

$$G - Z = G \cap \tilde{G} = \tilde{G} - A^*.$$

By Lemma 3.2 (iii) and by (3), the union of the sets of paths in (5) for all  $\mu < \mu^*$  is a set of disjoint paths. Thus, for  $U := \bigcup_{\mu < \mu^*} U_\mu$

*there is a set of disjoint  $U$ - $A$  paths whose set of first points is  $U$ ,  
and whose paths meet  $\tilde{G}$  only in  $U$ .* (6)

An important property of  $\tilde{G}$  is that the ends of  $G$  in  $B \cap \Omega$  correspond closely to ends of  $\tilde{G}$ . To establish this correspondence formally, we begin with the following observation:

*Every ray of an end  $\beta \in B$  has a tail in  $G - Z$ .* (7)

To see this, recall that all the  $D \in \mathcal{D}_1 \cup \mathcal{D}_2$  have pairwise disjoint closures, and that each of them is a connected subgraph of  $G$  whose closure contains an end or a vertex of  $A$ . Hence, a ray  $R$  of  $\beta$  meets only finitely many  $D \in \mathcal{D}_1 \cup \mathcal{D}_2$ , as we could otherwise find infinitely many disjoint  $R$ - $A$  paths, giving  $\bar{A} \cap B \neq \emptyset$  by Lemma 2.1 – a contradiction. Also,  $R$  meets every  $D \in \mathcal{D}_1 \cup \mathcal{D}_2$  only finitely often. Indeed,  $D$  lies in  $G_\mu$  for some  $\mu < \mu^*$  and is thus, by Lemma 3.2 (i), separated from  $\beta$  by its finite set of neighbours  $N(D)$ . This establishes (7).

*Let  $R_1, R_2$  be two rays in  $G \cap \tilde{G}$ , and assume that the end of  $R_1$  lies in  $B$ . Then  $R_1$  and  $R_2$  are equivalent in  $G$  if and only if they are equivalent in  $\tilde{G}$ .* (8)

To prove (8), suppose first that  $R_1, R_2$  are equivalent in  $G$ , i.e. belong to the same end  $\beta \in B$ . Then there is a ray  $R_3$  that meets both  $R_1$  and  $R_2$  infinitely often, and hence ends in  $\beta$ . By (7),  $R_3$  has a tail in  $G - Z = \tilde{G} - A^*$ , showing that  $R_1$  and  $R_2$  are equivalent also in  $\tilde{G}$ .

Conversely, if  $R_1$  and  $R_2$  are joined in  $\tilde{G}$  by infinitely many disjoint paths, we can replace any vertices  $a_\alpha \in \tilde{V} \setminus V = A^*$  on these paths by finite paths in  $D_\alpha$  to obtain infinitely many disjoint  $R_1$ - $R_2$  paths in  $G$ . This completes the proof of (8).

We can now define our correspondence between the ends in  $B$  and certain ends of  $\tilde{G}$ . For every end  $\beta \in B$  there is by (7) an end  $\beta' \in \tilde{\Omega}$  such that  $\beta \cap \beta' \neq \emptyset$ . By (8), this end  $\beta'$  is unique and the map  $\beta \mapsto \beta'$  is injective. Moreover,

$$\tilde{B} := (B \cap V) \cup \{\beta' \mid \beta \in B \cap \Omega\} \subseteq \tilde{V} \cup \tilde{\Omega}$$

by Lemma 3.2 (i). For each  $\mu < \mu^*$ , let

$$\tilde{A}_\mu := U_\mu \cup \{a_\alpha \mid \alpha \in O_\mu\},$$

if  $G_\mu \neq \emptyset$ ; if  $G_\mu = \emptyset$ , put  $A_\mu, \tilde{A}_\mu := \emptyset$ . Then let

$$\tilde{A} := \left( A \setminus \left( \bigcup_{\mu < \mu^*} A_\mu \cup I \right) \right) \cup \bigcup_{\mu < \mu^*} \tilde{A}_\mu,$$

which is a subset of  $\tilde{V}$  by Lemma 3.2 (iii),(iv). Finally, let

$$G' := \tilde{G} - (\tilde{A} \cap \tilde{B}).$$

To show the assertions (i)–(iv), we will apply Lemma 3.1 to the graph  $\tilde{G}$  and the sets  $\tilde{A}$  and  $\tilde{B}$ .

So, let us show that

$$(\overline{\tilde{A} \setminus \tilde{A}}) \cap \tilde{B} = \emptyset = \tilde{A} \cap (\overline{\tilde{B} \setminus \tilde{B}})$$

(with closures taken in  $|\tilde{G}|$ ). We trivially have  $\tilde{A} \cap (\overline{\tilde{B} \setminus \tilde{B}}) = \emptyset$  because  $\tilde{A} \subseteq \tilde{V}$ . To prove that  $(\overline{\tilde{A} \setminus \tilde{A}}) \cap \tilde{B} = \emptyset$ , consider an end  $\beta' \in \tilde{B}$ . The corresponding end  $\beta \in B$  has a neighbourhood  $C := \hat{C}(S, \beta)$  in  $|G|$  that avoids  $A$ . By (6), and since  $S$  is finite, the intersection  $C \cap U =: U_C$  is finite. Also, as in the proof of (7),  $C$  may meet only finitely many  $D_\alpha \in \mathcal{D}_2$ . Denote by  $O_C$  the set of the corresponding  $a_\alpha \in \tilde{G}$ . Adding to  $S \setminus Z$  the sets  $U_C$  and  $O_C$  then yields a finite set  $S' \subseteq \tilde{V}$  such that the neighbourhood  $\hat{C}'(S', \beta')$  in  $|\tilde{G}|$  even avoids  $\tilde{A}$ .

Thus, Lemma 3.1 is applicable and yields sets  $A' \subseteq V'$  and  $B' \subseteq V' \cup \Omega'$  satisfying (iii). For (i) use Lemma 3.1 (i), and observe that  $|\tilde{A}| \leq |A|$ . (Indeed,  $\tilde{A}$  is comprised of two sets, one of which is contained in  $A$ . The other set,  $\bigcup_{\mu < \mu^*} \tilde{A}_\mu$ , has cardinality at most  $|A|$  by (6).) Assertion (ii) follows from the definition of  $\tilde{B}$  and Lemma 3.1 (ii).

We now prove assertion (iv) of the lemma. Suppose  $G'$  satisfies the Erdős-Menger conjecture for  $A'$  and  $B'$ . Then, by Lemma 3.1, there is also in  $\tilde{G}$  a set  $\tilde{\mathcal{P}}$  of disjoint  $\tilde{A}$ – $\tilde{B}$  paths and an  $\tilde{A}$ – $\tilde{B}$  separator  $\tilde{X}$  on  $\tilde{\mathcal{P}}$ . In order to turn  $\tilde{\mathcal{P}}$  into a set  $\mathcal{P} := \{P \mid \tilde{P} \in \tilde{\mathcal{P}}\}$  of disjoint  $A$ – $B$  paths in  $G$ , consider any  $\tilde{P} \in \tilde{\mathcal{P}}$ . If the first point  $a$  of  $\tilde{P}$  lies in  $A$  we leave  $\tilde{P}$  unchanged, i.e. set  $P := \tilde{P}$ . If  $a \in \tilde{A} \setminus (A \cup A^*)$ , then  $a \in U_\mu$  for some  $\mu < \mu^*$ , and we let  $P$  be the union of  $\tilde{P}$  with an  $A_\mu$ – $U_\mu$  path in  $G_\mu$  that ends in  $a$ ; this can be done disjointly for different  $\tilde{P} \in \tilde{\mathcal{P}}$  if we use the paths from (6). Moreover, the  $A_\mu$ – $H_\mu$  path concatenated with  $\tilde{P}$  in this way has only its last vertex in  $\tilde{G}$ , so it will not meet any other vertices on  $\tilde{\mathcal{P}}$ . Finally if  $a = a_\alpha \in A^*$ , we let  $P$  be obtained from  $\tilde{P}$  by replacing  $a$  with a path in  $D_\alpha$  that starts at the end  $\alpha$  and ends at the vertex of  $D_\alpha$  incident with the first edge of  $\tilde{P}$  (the edge incident with  $a$ ). In all these cases we have  $P \subseteq G$ , because  $\tilde{P}$  has no vertex in  $A^*$  other than possibly  $a$ . And no vertex of  $P$  other than possibly its last vertex lies in  $B$ , because  $B \cap V = \tilde{B} \cap \tilde{V}$  and any new initial segment of  $P$  lies in a subgraph  $G_\lambda - S_\lambda$  of  $G$  which avoids  $B$  by Lemma 3.2 (i).

It remains to check that the paths  $P$  just defined have distinct last points in  $B$  even when the last points of the corresponding paths  $\tilde{P}$  are ends. However if  $\tilde{P}$  ends in  $\beta' \in \tilde{B}$  then its tail  $\tilde{P} - a \subseteq P \subseteq G$  is equivalent in  $\tilde{G}$  to some ray in  $\beta' \cap \beta$ , by definition of  $\beta'$ . By (8) this implies  $\tilde{P} - a \in \beta$ , so the last point of  $P$  is  $\beta \in B$ . And since the map  $\beta \mapsto \beta'$  is well defined, these last points differ for distinct  $P$ , because the corresponding paths  $\tilde{P}$  have different endpoints  $\beta'$  by assumption.

We still need an  $A$ - $B$  separator on  $\mathcal{P}$ . The only vertices  $x \in \tilde{X}$  that do not lie on the path  $P$  obtained from the path  $\tilde{P}$  containing  $x$  are points in  $A^*$ . So let  $X$  be obtained from  $\tilde{X}$  by replacing every end  $\beta' \in \tilde{X} \cap \tilde{B}$  with the corresponding end  $\beta \in B$  and replacing every  $a_\alpha \in \tilde{X} \cap A^*$  with the end  $\alpha \in A$ . Since  $P \in \mathcal{P}$  starts in  $\alpha$  if  $\tilde{P}$  starts in  $a_\alpha$  (and  $P$  ends in  $\beta$  if  $\tilde{P}$  ends in  $\beta'$ ), this set  $X$  consists of a choice of one point from every path in  $\mathcal{P}$ .

Let us then show that

$$X \text{ is an } A\text{-}B \text{ separator in } G. \quad (9)$$

Suppose there exists a path  $Q \subseteq G - X$  that starts in  $A$  and ends in  $B$ . Lemma 3.2 (v) enables us to choose  $Q$  as a path starting in  $A \setminus I$ . Our aim is to turn  $Q$  into an  $\tilde{A}$ - $\tilde{B}$  path  $Q'$  in  $\tilde{G}$  that avoids  $\tilde{X}$ , which contradicts the choice of  $\tilde{X}$ .

If  $Q$  meets  $\bigcup \mathcal{D}_1$ , it has a last vertex there by (7), in  $D \in \mathcal{D}_1(\lambda)$ , say. Its next vertex  $a$  lies in  $U_\lambda$ , by the definition of  $D$ . We then define (for the time being)  $Q'$  as the final segment  $aQ$  of  $Q$  starting at  $a$ . If  $Q$  has no vertex in  $\bigcup \mathcal{D}_1$ , then either the first point of  $Q$  is a vertex  $a \in A \cap \tilde{A}$  (in which case we put  $Q' := Q$ ), or  $Q$  starts at an end  $\alpha \in A \setminus I$ . By Lemma 3.2 (iv), there exists a  $\lambda < \mu^*$  such that  $\alpha \in \overline{G}_\lambda$ , which implies  $\alpha \in O_\lambda$ . We make  $a := a_\alpha$  the starting vertex of  $Q'$  and continue  $Q'$  along  $Q$ , beginning with the last  $D_\alpha$ - $\tilde{G}$  edge on  $Q$ . Our assumption of  $\alpha \notin X$  implies that  $a_\alpha \notin \tilde{X}$ , by the definition of  $X$ . Thus in the first two cases,  $Q'$  is now a path in  $G - \bigcup \mathcal{D}_1$ ; in the third,  $Q'$  is a path in  $(G - \bigcup \mathcal{D}_1)/D_\alpha$ , which starts at the vertex  $a \in \tilde{A}$  and avoids  $\tilde{X}$ .

However,  $Q'$  may still meet  $\mathcal{D}_2$ . And although we know from (7) that  $Q'$  has a last vertex in  $\bigcup \mathcal{D}_2$ , say in  $D_{\alpha'}$ , we cannot simply shorten  $Q'$  to a path  $a_{\alpha'}Q'$  in  $\tilde{G}$ , because it may happen that  $a_{\alpha'} \in \tilde{X}$ . Instead, we will use Lemma 3.4 to replace any segments of  $Q'$  that meet some  $D_\alpha \in \mathcal{D}_2$  (with  $a_\alpha \neq a$ ) by paths through the corresponding  $G_\mu$  that avoid  $\tilde{X}$ . As we only have to deal with a finite initial segment of  $Q'$  and the  $D_\alpha$  are all disjoint, we are able to modify  $Q'$  step by step. Eventually, we will obtain a (walk that can be pruned to a) path  $Q'$  in  $\tilde{G}$  that avoids  $\tilde{X}$ , yielding the desired contradiction.

So consider a segment of  $Q'$  that meets some  $D_\alpha \in \mathcal{D}_2$ . By definition of  $D_\alpha$  we may assume that segment to be a  $T_\mu$ -path  $sQ't$  in  $H_\mu$ , where  $\mu$  is such that  $D_\alpha \subseteq G_\mu$ . By definition of  $H'_\mu$  (which is a subgraph of  $\tilde{G}$  by Lemma 3.2 (iii), i.e. no parts of  $H'_\mu$  were deleted or contracted when we defined  $\tilde{G}$ ), there are  $|S_\mu| + 1$  paths from  $s$  to  $t$  in  $H'_\mu$  that are disjoint outside  $sQ't$ . But  $H'_\mu$  contains at most  $|S_\mu|$  vertices from  $\tilde{X}$ : since these lie on disjoint paths ending in  $\tilde{B}$  and  $S_\mu$  separates  $H'_\mu \subseteq G_\mu$  from  $B$  in  $G$  and hence from  $\tilde{B}$  in  $\tilde{G}$ , all of these paths must meet  $S_\mu$ . So one of our  $|S_\mu| + 1$   $s$ - $t$  paths in  $H'_\mu$  avoids  $\tilde{X}$ , and we can use this path to replace  $sQ't$  on  $Q'$ . This completes the proof of (9).  $\square$

We can now repeat the reduction for the ends of  $B$ .

**Lemma 3.6.** *Let  $G = (V, E, \Omega)$  be a graph, and let  $A, B \subseteq V \cup \Omega$  be such that  $A \cap \overline{B} = \emptyset = \overline{A} \cap B$ . Then there are a minor  $G'$  of  $G$  and sets  $A', B' \subseteq V(G')$  satisfying*

- (i)  $|A'| \leq |A|$ ,
- (ii)  $G$  satisfies the Erdős-Menger conjecture for  $A$  and  $B$  if  $G'$  satisfies it for  $A'$  and  $B'$ .

*Proof.* Apply Lemma 3.5 twice, once for  $A$  and once for  $B$ .  $\square$

By Lemma 3.6, we already have the main case of Theorem 2.2:

**Proposition 3.7.** *Let  $G = (V, E, \Omega)$  be a graph, let  $A, B \subseteq V \cup \Omega$  be such that  $(\overline{A} \setminus A) \cap B = \emptyset = A \cap (\overline{B} \setminus B)$ , and let  $A$  be countable. Then  $G$  satisfies the Erdős-Menger conjecture for  $A$  and  $B$ .*

*Proof.* Use Lemmas 3.1 and 3.6, and Theorem 1.1.  $\square$

The next section will be spent on strengthening Proposition 3.7 to the full generality of Theorem 2.2.

## 4 Countable separators

In this section we prove Theorem 2.2. But first, let us establish a weaker version in which the separator  $X$  also satisfies  $A \cap (\overline{X} \setminus X) = \emptyset$ :

**Lemma 4.1.** *Let  $G = (V, E, \Omega)$  be a graph, let  $A, B \subseteq V \cup \Omega$  satisfy  $A \cap \overline{B} = \emptyset = \overline{A} \cap B$ , and suppose there exists a countable  $A$ - $B$  separator  $X \subseteq V \cup \Omega$  in  $G$  with  $A \cap (\overline{X} \setminus X) = \emptyset$ . Then  $G$  satisfies the Erdős-Menger conjecture for  $A$  and  $B$ .*

*Proof.* First, we find a countable  $A$ - $B$  separator  $Y$  and a set  $\mathcal{P}_Y$  of disjoint  $Y$ - $A$  paths satisfying

- (a) in every  $y \in Y$  there starts a path of  $\mathcal{P}_Y$ ;
- (b)  $Y \cap \Omega \subseteq A \cup B$ ;
- (c)  $(\overline{Y} \setminus Y) \cap B = \emptyset = Y \cap (\overline{B} \setminus B)$ .

We may clearly assume that  $X \cap \Omega \subseteq A \cup B$ . As  $\overline{A} \cap B = \emptyset$ , this implies that  $X \cap (\overline{A} \setminus A) = \emptyset$ . As also  $A \cap (\overline{X} \setminus X) = \emptyset$  and  $X$  is countable, we can use Proposition 3.7 to obtain a set  $\mathcal{P}_1$  of disjoint  $X$ - $A$  paths and an  $X$ - $A$  separator  $Y$  on  $\mathcal{P}_1$  in  $G$ .

We claim that  $Y$  together with the set

$$\mathcal{P}_Y := \{yP \mid P \in \mathcal{P}_1 \text{ and } y \in Y \cap P\}$$

of disjoint  $Y$ - $A$  paths is as desired.

Indeed,  $Y$  is countable because  $\mathcal{P}_1$  is. Further,  $Y$  is an  $A$ - $B$  separator in  $G$ : any path starting in  $B$  and ending in  $A$  meets  $X$  and thus has a subpath  $P$  starting in  $X$  and ending in  $A$ . But then  $P$  also meets  $Y$ . The conditions (a) and (b) are easily checked; for the latter recall that we assumed  $X \cap \Omega \subseteq A \cup B$ .

Next, we show (c). By (b), any end  $\alpha \in Y \cap (\overline{B} \setminus B)$  lies in  $A \cap \overline{B}$ , which is empty by assumption. Now consider an end  $\beta \in (\overline{Y} \setminus Y) \cap B$ . Every neighbourhood of  $\beta$  contains a point of  $Y$ ; it even contains infinitely many points of  $Y$ , as otherwise we could find a neighbourhood containing no point of  $Y$ . Choose a neighbourhood  $\hat{C}(S, \beta)$  of  $\beta$  that contains no point of  $A$  (which is possible, as  $\overline{A} \cap B = \emptyset$ ). The infinitely many points of  $Y$  lying in  $\hat{C}(S, \beta)$  are linked to  $A$  by disjoint paths in  $\mathcal{P}_Y$ . All of these infinitely many paths must meet the finite set  $S$ , a contradiction. Therefore, (c) is proved.

Having found  $Y$  and  $\mathcal{P}_Y$ , we now apply Proposition 3.7 again, this time for the sets  $Y$  and  $B$ . Thus, we get a system  $\mathcal{P}_2$  of disjoint  $B$ - $Y$  paths and a  $B$ - $Y$  separator  $Z$  on  $\mathcal{P}_2$  in  $G$ .

Finally, let  $\mathcal{P}$  be the system of disjoint  $B$ - $A$  paths obtained by concatenating every  $B$ - $Y$  path  $P \in \mathcal{P}_2$  with the unique  $Y$ - $A$  path  $P' \in \mathcal{P}_Y$  from (a) that starts at the endpoint of  $P$ . Note that these are indeed paths: if  $P$  terminates in an end  $\omega \in Y$  then, by (b),  $\omega \in A$  or  $\omega \in B$ . In the former case  $P'$  is trivial, in the latter case  $P$ . Clearly,  $Z$  lies on  $\mathcal{P}$ . All that is left to show is that  $Z$  is also an  $A$ - $B$  separator. So consider a path  $P$  that starts in  $A$  and ends in  $B$ . By definition of  $Y$ ,  $P$  meets  $Y$  and thus has a subpath starting in  $Y$  and ending in  $B$ . This subpath cannot avoid the  $Y$ - $B$  separator  $Z$ .  $\square$

We can now complete the proof of Theorem 2.2, which we restate:

**Theorem 2.2.** *Let  $G = (V, E, \Omega)$  be a graph, let  $A, B \subseteq V \cup \Omega$  satisfy  $A \cap \overline{B} = \emptyset = \overline{A} \cap B$ , and suppose that there exists a countable  $A$ - $B$  separator  $X \subseteq V \cup \Omega$  in  $G$ . Then  $G$  satisfies the Erdős-Menger conjecture for  $A$  and  $B$ .*

*Proof.* We will start by constructing a set  $A' \subseteq A$  and a countable  $A'$ - $B$  separator  $X'$  with

$$(\overline{X'} \setminus X') \cap A' = \emptyset, \quad (10)$$

to which we will then apply Lemma 4.1.

Using Lemma 3.2 with

$$A_\Omega := (\overline{X} \setminus X) \cap A,$$

we find  $I, \mu^*$  and families  $(G_\mu)_{\mu < \mu^*}$  and  $(S_\mu)_{\mu < \mu^*}$  satisfying the assertions (i)–(v) of Lemma 3.2. As before, write  $C_\mu := G_\mu - S_\mu$ . Setting  $A' := A \setminus I$  and

$$X' := \bigcup_{\mu < \mu^*} S_\mu \cup \left( X \setminus \bigcup_{\mu < \mu^*} \overline{C_\mu} \right)$$

we claim that  $X'$  is a countable  $A'$ - $B$  separator satisfying (10).

To see that  $X'$  is countable, recall that the sets  $\overline{C_\mu}$  are disjoint for different  $\mu$  (Lemma 3.2 (iii)), and that each  $\overline{C_\mu}$  with  $S_\mu \neq \emptyset$  contains an end  $\alpha \in A_\Omega$  (Lemma 3.2 (ii)). Since  $\alpha \in \overline{X}$  (by definition of  $A_\Omega$ ), we have  $\overline{C_\mu} \cap X \neq \emptyset$ , so by the countability of  $X$  there are only countably many such  $\mu$ .

Let us now show that  $X'$  is an  $A'$ - $B$  separator. Consider a path  $Q$  from  $A'$  to  $B$ . Since  $X$  is an  $A$ - $B$  separator,  $Q$  must meet  $X$ . If  $Q$  meets  $X$  outside  $X'$ , it meets  $X$  in  $\overline{C_\mu}$  for some  $\mu < \mu^*$ . By Lemma 3.2 (i), however,  $Q$  cannot be contained in  $\overline{C_\mu}$ , so  $Q$  meets  $S_\mu \subseteq X'$ .

To prove (10), suppose there is an end  $\omega \in (\overline{X'} \setminus X') \cap A'$ . Let us show first that

$$\omega \notin \overline{X}. \quad (11)$$

Suppose otherwise; then we have either  $\omega \in (\overline{X} \setminus X) \cap A = A_\Omega$  or  $\omega \in X \setminus X'$ . In both cases (in the first by Lemma 3.2 (iv), observe that  $\omega \notin I$  as  $\omega \in A'$ ; in the second by construction of  $X'$ ) there is a  $\mu < \mu^*$  with  $\omega \in \overline{G_\mu}$ . Then by Lemma 3.2 (iii), the sets of the form  $\hat{C}(S_\mu, \omega)$  are neighbourhoods of  $\omega$  that avoid  $X'$ , contradicting  $\omega \in \overline{X'} \setminus X'$ .

By (11), there exists a finite set  $S$  such that  $C(S, \omega) \cap X = \emptyset$ . By our choice of  $\omega$ , however,  $C(S, \omega)$  contains infinitely many vertices from  $X'$ , and hence from  $X' \setminus X \subseteq \bigcup_{\mu < \mu^*} S_\mu$ . Since each  $S_\mu$  is finite, we can thus find infinitely many  $\mu < \mu^*$  and corresponding vertices  $s_\mu \in S_\mu \cap C(S, \omega)$  that are distinct for different  $\mu$ . By Lemma 3.2 (ii) and  $S_\mu = N(C_\mu)$ , each  $s_\mu$  sends a path  $P_\mu \subseteq \overline{G[C_\mu \cup \{s_\mu\}]}$  to an end in  $A_\Omega \subseteq \overline{X}$ . These  $P_\mu$  are disjoint by Lemma 3.2 (iii), so only finitely many of them meet  $S$ . Every other  $P_\mu$  lies entirely in  $\overline{C(S, \omega)}$ , so  $\overline{C(S, \omega)} \cap \overline{X} \neq \emptyset$ . But then also  $C(S, \omega) \cap X \neq \emptyset$ , contradicting our choice of  $S$ . This establishes (10).

Applying Lemma 4.1 to  $A', B$  and the separator  $X'$  (note that  $\overline{A'} \cap B = \emptyset = A' \cap \overline{B}$ , as  $A' \subseteq A$ ), we obtain a set  $\mathcal{P}$  of disjoint  $A'$ - $B$  paths and an  $A'$ - $B$  separator on  $\mathcal{P}$ , which by Lemma 3.2 (v) is also an  $A$ - $B$  separator.  $\square$

## 5 Disjoint closures

We restate our second main result, which we shall prove in this section.

**Theorem 2.3.** *Every graph  $G = (V, E, \Omega)$  satisfies the Erdős-Menger conjecture for all sets  $A, B \subseteq V \cup \Omega$  that have disjoint closures in  $|G|$ .*

Our proof follows that of Theorem 1.2 as given in [7], and in addition we will draw on techniques from the proof of Lemma 3.5. Our aim is to reduce our problem to rayless graphs, and then apply Theorem 2.4. We thus need to dispose of the ends in  $G$ , which will be achieved in three steps. First, we delete all ends in  $\overline{A}$ . More precisely, we reduce the problem to a minor  $G'$  of  $G$  and to sets  $A', B'$  so that the closure of  $A'$  contains no ends. In the next step we repeat this procedure for  $B'$ . To preserve what we have gained in the first step, we have to be careful that no new ends are introduced into  $\overline{A'}$ . All this amounts to the following lemma:

**Lemma 5.1.** *Let  $G = (V, E, \Omega)$  be a graph, and let  $A, B \subseteq V \cup \Omega$  be such that  $\overline{A} \cap \overline{B} = \emptyset$ . Then there exists a minor  $G' = (V', E', \Omega')$  of  $G$  and sets  $A', B' \subseteq V' \cup \Omega'$  that satisfy the following conditions:*

- (a)  $\Omega' \cap \overline{A'} = \emptyset$  (in particular  $A' \subseteq V'$ );
- (b) if  $\Omega \cap \overline{B} = \emptyset$  then  $\Omega' \cap \overline{B'} = \emptyset$  (and in particular  $B' \subseteq V'$ );
- (c)  $\overline{A'} \cap \overline{B'} = \emptyset$ ; and
- (d) the Erdős-Menger conjecture holds for  $A$  and  $B$  in  $G$  if it holds for  $A'$  and  $B'$  in  $G'$ .

Two applications of Lemma 5.1, one for  $A$  and another for  $B$ , reduces our problem to the case that  $G$  has no ends in  $\overline{A} \cup \overline{B}$ . We then eliminate the remaining ends by the following lemma from [7]:

**Lemma 5.2.** [7] *Let  $G = (V, E, \Omega)$  be a graph, and let  $A, B \subseteq V$  be such that  $\Omega \cap (\overline{A} \cup \overline{B}) = \emptyset$ . Then  $G$  has a rayless subgraph  $G' \subseteq G$  containing  $A \cup B$  such that the Erdős-Menger conjecture for  $A$  and  $B$  holds in  $G$  if it does in  $G'$ .*

Finally, we apply Theorem 2.4 to complete the proof.

It thus remains to establish Lemma 5.1. We shall need the following easy lemma from [9]:

**Lemma 5.3.** [9] *Let  $G$  be a connected graph, and let  $U \subseteq V(G)$  be an infinite set of vertices. Then  $G$  contains either a comb with  $|U|$  teeth in  $U$  or a subdivided star with  $|U|$  leaves in  $U$ .*

*Proof of Lemma 5.1.* First, we construct a subgraph  $M \subseteq G$  whose closure does not contain ends of  $\bar{A}$ . More formally, our aim is that in  $|G|$

$$\Omega \cap \bar{A} \cap \bar{M} = \emptyset. \quad (12)$$

Our desired graph  $G'$  will then be obtained from a supergraph of  $M$ .

We define  $M$  by transfinite ordinal recursion, as a limit  $M = \bigcap_{\mu \leq \mu^*} M_\mu$  of a well-ordered descending family of subgraphs  $M_\mu$  indexed by ordinals. Put  $M_0 := G$ , and for a limit ordinal  $\mu \neq 0$  let  $M_\mu := \bigcap_{\mu' < \mu} M_{\mu'}$ . Now, consider a successor ordinal  $\mu + 1$ . If  $\Omega \cap \bar{A} \cap \bar{M}_\mu = \emptyset$  put  $\mu^* := \mu$  and  $M := M_\mu$ , and terminate the recursion. Otherwise, there is an  $\alpha_\mu \in \Omega \cap \bar{A} \cap \bar{M}_\mu$ . Since  $\bar{A} \cap \bar{B} = \emptyset$ , we can choose a finite vertex set  $L_\mu$  such that the open neighbourhood  $\hat{C}(L_\mu, \alpha_\mu)$  is disjoint from  $B$ . Put  $C_\mu := C(L_\mu, \alpha_\mu)$  and  $M_{\mu+1} := M_\mu - C_\mu$ . Observe that  $C_\mu \cap M_\mu$  is never empty as  $\alpha_\mu \in \bar{M}_\mu$ . Thus, the recursion terminates.

Let  $\mathcal{C}$  be the set of components of  $G - M$ . For every  $C \in \mathcal{C}$  put  $S_C := N_G(C)$ ,  $G_C := G[V(C) \cup S_C]$  and  $A_C := A \cap \bar{G}_C$ . We shall now proceed in a similar way as in the proof of Lemma 3.5: in order to obtain  $G'$  we will delete and contract certain connected subgraphs, each of which will lie in a  $C \in \mathcal{C}$ . Thus, the  $G_C$  here play a similar role as the  $G_\mu$  in Lemma 3.5. Now, in contrast to the set of neighbours  $S_\mu$  we deal with there, the sets  $S_C$  here are not necessarily finite. Consequently, it may happen that two  $C \in \mathcal{C}$  do not have disjoint closures. However, for our purposes it is sufficient to know that no end in  $A$  lies in the closure of two elements of  $\mathcal{C}$ :

$$A_C \cap A_D \subseteq S_C \cap S_D \text{ for two different } C, D \in \mathcal{C} \quad (13)$$

Indeed, suppose there is an end  $\alpha \in A_C \cap A_D = A \cap \bar{G}_C \cap \bar{G}_D$ . Then every neighbourhood  $\hat{C}(T, \alpha)$  of  $\alpha$  contains vertices of both  $G_C$  and  $G_D$ . Being connected,  $C(T, \alpha)$  also contains a  $G_C - G_D$  path, and therefore a vertex of  $S_C \subseteq V(M)$ . This implies  $\alpha \in \bar{M}$ , a contradiction to (12). As no vertices other than those of  $S_C \cap S_D$  lie in  $A_C \cap A_D$ , (13) is proved.

Now, consider a given  $C \in \mathcal{C}$ . Although  $S_C$  may be infinite it can be separated from  $A_C$  by finitely many vertices and ends. Indeed, we claim that

$$\bar{G}_C \text{ contains a finite set } \mathcal{P}_C \text{ of disjoint } S_C - A_C \text{ paths and a } S_C - A_C \text{ separator } X_C \text{ on } \mathcal{P}_C. \quad (14)$$

First, consider a non-trivial  $S_C - A_C$  path  $P$ .  $P$  is either completely contained in  $G_C$  or in  $G - C$ . Suppose the latter. Then  $P$  terminates in an end  $\alpha \in A_C$ . By (13),  $P$  can meet every  $D \in \mathcal{C}$  only finitely often, hence it meets  $M$  infinitely often, a contradiction to (12). Thus, every  $S_C - A_C$  path lies completely in  $G_C$ .

Next, let us show that each set of disjoint  $S_C - A_C$  paths in  $G_C$  is finite. So suppose there is an infinite set  $P_1, P_2, \dots$  of such paths. Let  $S \subseteq S_C$  be the set of first vertices of the  $P_i$ . We claim there is a comb in  $G$  with teeth in  $S$ . If not, applying Lemma 5.3 to  $H := C \cup \bigcup_{i \in \mathbb{N}} P_i$  we obtain a subdivided infinite star with leaves in  $S$ . Since each vertex  $s \in S$  has degree 1 in  $H$ , the centre  $v$  of the star must lie in  $C$ . Let  $\mu$  be the step when  $v$  was deleted from  $G$ , i.e.

$\mu := \min\{\mu' \mid v \in V(C_{\mu'})\}$ . Then, the finite set  $L_\mu$  separates  $v$  from  $S$ , which is impossible. Thus, there is a comb with teeth in  $S$ . Let  $\omega \in \Omega$  be the end of its spine. Then every neighbourhood of  $\omega$  contains infinitely many vertices of  $S$  and then also infinitely many of the  $P_i$ . Consequently, infinitely many elements of  $A$  lie in the neighbourhood. Thus,  $\omega \in \overline{S_C} \cap \overline{A} \subseteq \overline{M} \cap \overline{A}$ , a contradiction to (12).

Taking an inclusion-maximal (and hence finite) set of disjoint  $S_C$ - $A_C$  paths in  $G$  we see that  $S_C$  is separated from  $A_C$  in  $G$  by a countable separator (namely by the union of the vertex sets of the paths together with the set of last points). Furthermore, (12) implies  $\overline{S_C} \cap A_C \cap \Omega = \emptyset$ . Thus, by Lemma 3.1 and Theorem 2.2, there is a set  $\mathcal{P}_C$  of disjoint  $S_C$ - $A_C$  paths in  $G$  and a separator  $X_C$  on  $\mathcal{P}_C$ . By the preceding arguments,  $\mathcal{P}_C$  is a finite set of paths in  $G_C$ , which establishes (14).

For  $C \in \mathcal{C}$ , put  $U_C := X_C \cap V$  and  $O_C := X_C \cap \Omega$ . As in the proof of Lemma 3.5 we will delete certain subgraphs of  $G_C$  that lie “behind”  $U_C$  and contract others that contain ends of  $O_C$ . First of all, we see that  $U_C$  separates  $S_C$  from  $A_C \setminus O_C$ . This is exactly the same as (3) in the proof of Lemma 3.5.

We then define  $\mathcal{D}_1(C)$  as the set of all the components  $D$  of  $G - U_C$  whose closure  $\overline{D}$  meets  $A_C \setminus O_C$ . Because  $U_C$  is an  $S_C$ - $(A_C \setminus O_C)$  separator, these components satisfy  $D \subseteq G_C - S_C$ , and their neighbourhood  $N(D) \subseteq U_C$  in  $G$  is finite. Also,  $\overline{D} \cap O_C = \emptyset$  for all  $D \in \mathcal{D}_1(C)$ , which can be proved as (4) of Lemma 3.5. We put  $H_C := G_C - \bigcup \mathcal{D}_1(C)$  and see that  $G_C$  contains a set of disjoint  $H_C$ - $A_C$  paths whose set of first points is  $U_C$ .

To find the subgraphs which will be contracted, we extend in the proof of Lemma 3.5 the finite set  $U_\mu \cup S_\mu$  to a finite set  $T_\mu$  that separates the ends in  $O_\mu$  pairwise. This enables us to apply Lemma 3.4. Here,  $S_C$  may be an infinite set, so to find a suitable finite set  $T_C$  we cannot use a superset of  $S_C$ . Instead, consider for each  $\alpha \in O_C$  the ordinal  $\mu_\alpha$  for which  $\alpha \in C_{\mu_\alpha}$ . Then, the finite set  $L_{\mu_\alpha}$  separates  $\alpha$  from  $S_C \subseteq V(M)$ . The union of these finitely many finite sets together with  $U_C$  can be taken as the finite set  $T_C$  we need for Lemma 3.4.

Define  $H'_C$  as the finite subgraph of  $H_C$  containing  $T_C$  which Lemma 3.4 provides for  $k := |T_C| + 1$ , and for each  $\alpha \in O_C$  let  $D_\alpha$  be the component of  $G - H'_C$  to which  $\alpha$  belongs. Finally, we set  $\mathcal{D}_2(C) := \{D_\alpha \mid \alpha \in O_C\}$  and  $\mathcal{D}_i := \bigcup_{C \in \mathcal{C}} \mathcal{D}_i(C)$ ,  $i = 1, 2$ . By deleting all subgraphs in  $\mathcal{D}_1$  and by contracting each subgraph  $D_\alpha \in \mathcal{D}_2$  to a vertex  $a_\alpha$  we obtain the graph  $\tilde{G} = (\tilde{V}, \tilde{E}, \tilde{\Omega})$ . We put

$$\tilde{A} := (A \setminus \bigcup_{C \in \mathcal{C}} A_C) \cup \bigcup_{C \in \mathcal{C}} U_C \cup \{a_\alpha \mid D_\alpha \in \mathcal{D}_2\}.$$

Observe that  $\tilde{A} \subseteq \tilde{V}$  because of (12).

Let us show that (a) holds for  $\tilde{A}$  in  $\tilde{G}$ . Indeed, suppose not. Then, there is by Lemma 2.1 a comb  $\tilde{K} \subseteq \tilde{G}$  with teeth in  $\tilde{A}$  and spine  $R$ , say. Consider a  $R$ - $\tilde{A}$  path  $P$  of the comb. If  $P$  ends in a vertex of  $U_C$  for some  $C \in \mathcal{C}$  then we extend  $P$  using the corresponding path in  $\mathcal{P}_C$  to a  $R$ - $A$  path. If  $P$  ends in one of the contracted vertices  $a_\alpha$  we substitute its last edge by a path through  $D_\alpha$  ending in  $\alpha$  so that again we obtain a  $R$ - $A$  path. Because of (13), and since the  $D_\alpha$  and all the paths in  $\bigcup_{C \in \mathcal{C}} \mathcal{P}_C$  are disjoint, all these changed paths are still disjoint. Thus we have found a comb  $K \subseteq G$  with spine  $R$  and teeth in  $A$ . Only finitely many of the  $R$ - $A$  paths of  $K$  may meet  $M$  as otherwise the end of  $R$  lies in the closure of both  $A$  and  $M$ , a contradiction to (12). Thus, we may

assume that  $K$  is contained in a  $C \in \mathcal{C}$ . In particular,  $R \subseteq C \cap \tilde{G}$ . But the finite set  $X_C$  separates  $C \cap \tilde{G}$  from  $A_C$ , contradicting that  $K$  contains infinitely many disjoint  $R$ - $A_C$  paths. This proves (a) for  $\tilde{A}$ .

Making use of (13), we see in a similar way as in the proof of Lemma 3.5, that every ray  $R$  of an end  $\beta \in B$  has a tail in  $G \cap \tilde{G}$ . There thus is a  $\beta' \in \tilde{\Omega}$  with  $\beta \cap \beta' \neq \emptyset$  and it can be shown as in the proof of Lemma 3.5 that this mapping is injective. Put

$$\tilde{B} := (B \cap V) \cup \{\beta' \mid \beta \in B \cap \Omega\} \subseteq \tilde{V} \cup \tilde{\Omega}.$$

The graph  $\tilde{G}$  together with the sets  $\tilde{A}$  and  $\tilde{B}$  is almost what we want. Indeed, (a) is satisfied and below we shall see that (b) and (d) hold as well. Only (c) fails since  $\tilde{A}$  and  $\tilde{B}$  may share some vertices. However,  $\overline{\tilde{A}}$  does not contain ends by (a). Thus

$$\emptyset = \overline{\tilde{A}} \cap \overline{\tilde{B}} \cap \tilde{\Omega} = (\overline{\tilde{A}} \setminus \tilde{A}) \cap (\overline{\tilde{B}} \setminus \tilde{B}).$$

Now, in a similar way as in Lemma 3.1 we find sets  $A'$  and  $B'$  in the graph  $G' := \tilde{G} - (\tilde{A} \cap \tilde{B})$  such that (c) holds and properties (a), (b) and (d) are preserved. Therefore, to finish the proof it suffices to show (b) and (d) for  $\tilde{A}$  and  $\tilde{B}$  in  $\tilde{G}$ .

For (b), assume that  $\tilde{\Omega} \cap \overline{\tilde{B}} \neq \emptyset$ . Then either  $\tilde{B}$  contains ends, or there is an end  $\omega \in \tilde{\Omega}$  such that every neighbourhood  $\hat{C}(S, \omega)$  in  $\tilde{G}$  contains vertices of  $\tilde{B}$ . In the latter case, by (a), we may suppose that  $\hat{C}(S, \omega)$  does not contain vertices of  $\tilde{A}$ , and thus view it as a neighbourhood in  $G$  containing vertices of  $B$ . Thus, in both cases,  $\Omega \cap \overline{B} \neq \emptyset$ , which proves (b).

The proof of (d) is completely analogous to the proof of (iv) in Lemma 3.5. Note that when we change the set  $\tilde{\mathcal{P}}$  of disjoint  $\tilde{A}$ - $\tilde{B}$  paths into  $A$ - $B$  paths, the paths still have, by (13), distinct first points.  $\square$

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